# Nurse Staffing under Absenteeism: A Distributionally Robust Optimization Approach

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• Nurse staffing is <u>significant</u>.



Nurse staffing cost [1]

Safe nurse-to-patient ratio (NPR) [2]

[1] <u>https://marketrealist.com/2014/11/analyzing-hospital-expenses</u>
[2] <u>https://www.nationalnursesunited.org/ratios</u>

Four Phases of nurse planning [3]:

- 1. Demand forecast and staffing.
- 2. Nurse shift scheduling.
- 3. Pre-shift staffing and re-scheduling.
- 4. Nurse-patient assignment.

Our research:

- Phases 1 and 3.
- Outputs can be used in Phases 2 and 4.



[3] Jonathan F Bard. Nurse scheduling models. *Wiley Encyclopedia of Operations Research and Management Science*, 2010.

- Nurse staffing is <u>challenging</u>.
  - Uncertainties:
    - Demand side (nurse demand).

Random patient census.

- Exogenous.
- Supply side (nurse absenteeism).
  - # nurses showing up for the shift.
  - ✓ Work-related stress; "vicious cycle."



- Nurse staffing is <u>challenging</u>.
  - Uncertainties:
    - Demand side (nurse demand).
      - ✓ Random patient census.
      - Exogenous.
    - Supply side (nurse absenteeism).
      - # nurses showing up for the shift.
      - ✓ Work-related stress; "vicious cycle."
      - Depends on the <u>staffing level</u>. [4]
      - Endogenous.
  - Temporary nurses:
    - Called in to maintain the safe NPR.

[4] Green et al., "Nursevendor Problem": Personnel Staffing in the Presence of Endogenous Absenteeism. *Management Science*, 2013.

- Potential solution:
  - ▶ Float pools.
    - Pool : a group of units.
    - Pool structure : which units belong to which pools.



- Potential solution:
  - Pool nurses.
    - Assigned to a unit within the pool *after uncertainties are realized*.



- Increases operational flexibility.
- Reduce # temporary nurses.
- Higher quality of care.
- Lower nurse staffing cost.

- Potential solution:
  - Cross-training.
    - Pool nurses trained for skills required by units within the pool.
    - Cross-training nurses causes an extra cost.



- Minimize # cross-training by optimizing the pool structure.

# **Very Brief Literature Review**

- Majority: deterministic models:
  - Ignore demand and absenteeism uncertainties.
  - Underestimate the staffing costs. [5]
- Stochastic models:
  - Majority: demand uncertainty.
    - Ignore absenteeism; "results in under-staffing." [4]
  - Absenteeism leads to endogenous uncertainty.
    - Computationally prohibitive. [6]
    - Heuristics typically used. [7], [8]

[5] Kao and Queyranne, Budgeting Costs of Nursing in a Hospital. *Management Science*, 1985.
[6] Easton, Service Completion Estimates for Cross-trained Workforce Schedules under Uncertain Attendance and Demand. *Production and Operations Management*, 2014.
[7] Kayse L. Maass et al., Incorporating nurse absenteeism into staffing with demand uncertainty. *Health care management science*, 20(1), 2017.
[8] Wang and Gupta, Nurse Absenteeism and Staffing Strategies for Hospital Inpatient Units. Manufacturing & Operations Service Management, 2014.

# **Very Brief Literature Review**

- Challenge of modeling endogenous absenteeism.
  - Learn-and-optimize turns to under-staff. [4]



# **Very Brief Literature Review**

- Potential walk-around:
  - Parametric models. [4]
    - Independence, homogeneous absence probability.
    - Data-driven; closed-form solution for 1-unit staffing.
    - Extension to multiple-unit & multiple-pool?
    - What if the chosen model is **biased**?





# **Our Proposal**

- Non-Parametric Model:
  - A family of probability distributions.
  - Moment-based ambiguity set:
    - Moments of nurse demands (support, mean, variance, etc.).
    - # nurses who show up  $\leq$  staffing level.
    - Mean of # nurses who show up = f(staffing level).
    - Data-driven.
- Distributionally robust optimization:
  - Multiple-unit, multiple-pool.
  - MILP representable; global optimality.
  - Extension to optimal pool structure design.

### Outline

- Distributionally Robust Nurse Staffing (DRNS) Model:
  - A two-stage model with a moment-based ambiguity set.
  - Recast it as a deterministic two-stage min-max formulation.
  - Separation approach.
- More Tractable Reformulation:
  - Special pool structures.
    - (1) One pool, (2) Disjoint pools, (3) Chained pools.
  - A monolithic MILP reformulation.
- Optimal Nurse Pool Design (ONPD).
- Numerical Case Studies.

- Variables (1<sup>st</sup>-stage, here-and-now):
  - $w_j = \#$  nurses assigned to unit *j* (regular nurses).
  - $y_i = \#$  nurses assigned to pool *i* (pool nurses).
- Objective:
  - Minimize the worst-case expected staffing cost.



• Ambiguity set:

$$\mathcal{D} = \left\{ \mathbb{P} : \left( \mathbb{E}_{\mathbb{P}}[\tilde{\eta}_{j}^{q}] = \mu_{jq}, \quad \forall j \in [J], \; \forall q \in [Q], \right) \\ \left( \mathbb{E}_{\mathbb{P}}[\tilde{w}_{j}] = f_{j}(w_{j}), \quad \forall j \in [J] : w_{j} \geq 1, \right) \\ \mathbb{E}_{\mathbb{P}}[\tilde{y}_{i}] = g_{i}(y_{i}), \quad \forall i = [I] : y_{i} \geq 1 \right\}$$

- Nurse demand (*Exogenous*):
  - Marginal moment information.
- Nurse absenteeism (*Endogenous*):
  - Mean of # nurses who show up = f(staffing level).
  - *f*, *g*: Piecewise linear function in general.

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- Nurse demand (*Exogenous*):
  - Marginal moment information.
- Nurse absenteeism (*Endogenous*):
  - Mean of # nurses who show up = f(staffing level).
  - *f*, *g*: Piecewise linear function in general.
  - This talk: linear *f*, *g* for simplicity.

• Recourse problem (after uncertainties are realized):

$$\begin{split} V(\tilde{w}, \tilde{y}, \tilde{\eta}) &= \min_{x, z, e} \quad \sum_{j=1}^{J} \left( c^{\mathsf{x}} x_j - c^{\mathsf{e}} e_j \right) \\ \text{s.t.} \quad \tilde{w}_j + \sum_{i: j \in P_i} z_{ij} + x_j - e_j = \tilde{\eta}_j, \quad \forall j = 1, \dots, J, \\ \sum_{j \in P_i} z_{ij} = \tilde{y}_i, \quad \forall i = 1, \dots, I, \\ x_j, e_j \in \mathbb{Z}_+, \quad \forall j = 1, \dots, J, \\ z_{ij} \in \mathbb{Z}_+, \quad \forall i = 1, \dots, I, \quad \forall j \in P_i. \end{split}$$



• Recourse problem (after uncertainties are realized):  $V(\tilde{w}, \tilde{y}, \tilde{\eta}) = \min_{x, z, e} \sum_{j=1}^{J} (c^{x}x_{j} - c^{e}e_{j})$ s.t.  $\tilde{w}_{j} + \sum_{i:j \in P_{i}} z_{ij} + x_{j} - e_{j} = \tilde{\eta}_{j}, \quad \forall j = 1, \dots, J,$   $\sum_{j \in P_{i}} z_{ij} = \tilde{y}_{i}, \quad \forall i = 1, \dots, I,$   $x_{j}, e_{j} \in \mathbb{Z}_{+}, \quad \forall j = 1, \dots, J,$   $z_{ij} \in \mathbb{Z}_{+}, \quad \forall i = 1, \dots, I, \quad \forall j \in P_{i}.$ 

#### **Proposition.**

The IP recourse problem is equivalent to the following linear program:

$$V(\tilde{w}, \tilde{y}, \tilde{\eta}) = \max_{(\alpha, \beta) \in \Lambda} \sum_{j=1}^{J} (\tilde{\eta}_j - \tilde{w}_j) \alpha_j + \sum_{i=1}^{J} \tilde{y}_i \beta_i$$

where

$$\Lambda := \{ \alpha \in \mathbb{R}^J, \beta \in \mathbb{R}^I \ : \beta_i + \alpha_j \leq 0, \ \forall i \in [I], \ \forall j \in P_i, \ \alpha_j \in [c^{\text{e}}, c^{\text{x}}], \ \forall j \in [J] \}$$

#### **Proposition.**

For given w and y, the worst-case expectation  $\sup_{\mathbb{P}\in\mathcal{D}} \mathbb{E}_{\mathbb{P}}[V(\tilde{w}, \tilde{y}, \tilde{\eta})]$  is equal to the optimal objective value of the following min-max optimization problem:

$$\min_{\gamma \ge 0, \ \lambda \ge 0, \ \rho} \left[ \max_{(\alpha,\beta) \in \Lambda} F(\alpha,\beta) - \sum_{i=1}^{I} (c_i y_i + d_i) \lambda_i + \sum_{j=1}^{J} \left[ \sum_{q=1}^{Q} \mu_{jq} \rho_{jq} - (a_j w_j + b_j) \gamma_j \right]$$
  
s.t.  $\gamma_j \le M w_j, \ \forall j = 1, \dots, J,$   
 $\lambda_i \le M y_i, \ \forall i = 1, \dots, I,$   
where  $F(\alpha,\beta) := \sum_{j=1}^{J} \left[ (-\alpha_j + \gamma_j) w_j \right]^+ + \sum_{i=1}^{I} \left[ (\beta_i + \lambda_i) y_i \right]^+ + \sum_{j=1}^{J} \sup_{\tilde{\eta}_j \in [\underline{\eta}_j, \overline{\eta}_j]_{\mathbb{Z}}} \left\{ \alpha_j \tilde{\eta}_j - \sum_{q=1}^{Q} \rho_{jq} \tilde{\eta}_j^q \right\}$ 

- How do we solve the min-max problem?
  - Separation approach :  $V \ge F(\alpha, \beta), \forall (\alpha, \beta) \in \Lambda$
  - *Reformulate* the inner maximization problem :  $\max_{(\alpha,\beta)\in\Lambda} F(\alpha,\beta)$
  - Search the extreme point and direction of the polyhedron  $\Lambda$

#### **Proposition.**

For given  $w, y, \gamma_j, \lambda_i$ , and  $\rho$ , problem  $\max_{(\alpha,\beta)\in\Lambda} F(\alpha,\beta)$  has the same optimal objective value as the following integer (linear) program:

**IP Sub-Problem** 
$$\max_{t,s,r,p} \sum_{j=1}^{J} \left( c_j^{t} t_j + c_j^{r} r_j \right) + \sum_{i=1}^{I} \left( c_i^{p} p_i + \sum_{j \in P_i} c_i^{s} s_{ij} \right)$$
$$\text{s.t.} \quad (t,s) \in \mathcal{H}$$
$$t_j + r_j = 1, \quad \forall j = 1, \dots, J,$$
$$\sum_{j \in P_i} s_{ij} + p_i = 1, \quad \forall i = 1, \dots, I.$$

★ The DRNS model = Deterministic two-stage min-max problem.
★ Solved by the separation approach.

### **More Tractable Reformulation**

• Structure 1 : One pool



• Structure 2 : Disjoint pools



#### Theorem.

In structure 1 and 2,

#### LP relaxation of *H* = Convex hull of *H*

# **More Tractable Reformulation**

• Structure 3 : Chained pools



In structure 3,

#### LP relaxation of H $\neq$ convex hull of H

#### **Proposition.**

In structure 3, the IP sub-problem = two longest path problems.

★ The DRNS model under the one pool, disjoint pools, and chained pool structures can be expressed as *a monolithic MILP* which can be solved by the *branch-and-bound algorithm*.

# **ONPD** Model

- Motivation
  - Two pool structures, same cost:



• A hospital hiring pool nurses should have a pool structure that minimizes the number of cross-training and the staffing cost.

# **ONPD Model**

- The ONPD model is built upon the DRNS model for disjoint pools
  - Additional binary variables:

$$y_{ij} = \begin{cases} 1 \text{ if unit } j \text{ is in pool } i \\ 0 \text{ otherwise} \end{cases}$$

- Objective:
  - minimize the #cross-training.
- Additional constraints:
  - a target staffing cost (budget).
- ★ The ONPD model provides *an optimal disjoint pool structure* that minimizes the #cross-training while achieving a target staffing cost

#### **1. Computational performance**

Two-stage min-max problem (separation approach)

VS

A monolithic MILP (branch-and-bound)

- I = number of pools, J = number of units
- For each [I, J], we generate 10 instances to calculate the *average* computational time

One pool structure			Disjoint pools structure		
[I,J]	Two-stage	MILP	[I,J]	Two-stage	MILP
[1, 5]	2.28	0.09	[3,5]	1.93	0.09
[1,7]	8.24	0.10	[3,7]	9.82	0.10
[1, 10]	44.42	0.13	[3, 10]	68.33	0.16
[1, 20]	> 3600	0.37	[3, 20]	> 3600	0.37
[1, 50]	> 3600	1.51	[3, 50]	> 3600	1.13

Average CPU Seconds

#### 2. A small example

• 4 units with *one pool structure* 

Input parameters (DRNS model)						
Nurse staffing costs	Bounds on number of nurses	Nurse demands				
A regular nurse : \$400	Unit 1: $w_1 \in [0, 10]$	Unit 1: $\tilde{\eta}_1 \in [1, 30]$ , mean = 22.48				
A pool nurse : \$425	Unit 2: $w_2 \in [0, 9]$	Unit 2: $\tilde{\eta}_2 \in [1, 28]$ , mean = 21.70				
A temporary nurse : \$460	Unit 3: $w_3 \in [0, 11]$	Unit 3: $\tilde{\eta}_3 \in [1, 32]$ , mean = 25.18				
	Unit 4: $w_4 \in [0, 11]$	Unit 4: $\tilde{\eta}_4 \in [1, 31]$ , mean = 25.39				
	Pool: $y_1 \in [0, 10]$					

Output results (DRNS model)				
Worst-case expected cost	WE[cost] = 13468.2			
Nurse staffing decision	$w_1=6, w_2=6, w_3=7, w_4=7, y_1=0$			

Computational performances						
	Time [s]	Node [#]	Gap [%]			
Two-stage	1.49	2321	0			
MILP	0.02	0	0			

- GUROBI 7.0.1 (with Python 2.7)
- Processor: Intel® Core<sup>TM</sup>i7-4850 HQ CPU@2.30GHz / RAM: 16GB / OS: 64bit

#### 2. When should we construct float pools?

- Experiment
  - Inputs:

 $TCR = \frac{\text{cost of hiring a temporary nurse}}{\text{cost of hiring a regular nurse}}$  $PCR = \frac{\text{cost of hiring a pool nurse}}{\text{cost of hiring a regular nurse}}$  $ARR = \frac{\text{absenteeism rate of pool nurses}}{\text{absenteeism rate of regular nurses}}$ 

- Output:

$$OVG[\%] = \frac{WE[cost]_0 - WE[cost]_1}{WE[cost]_0} \times 100$$
$$= Value of Float Pools$$

- 2. When should we construct float pools?
  - Results:



- 2. When should we construct float pools?
  - Results:



#### 2. When should we construct float pools?

- Intuition: hiring pool nurses benefits when
  - 1. Cost of temporary nurses are high,
  - 2. Cost of pool nurses are low, and
  - 3. Pool nurses have low absenteeism.

#### 3. Why should we model absenteeism?

- Out-of-sample simulation
  - 1. Compute the staffing decision *considering / ignoring* absenteeism.
  - 2. Simulate the out-of-sample average costs.
- Result (E[cost]<sub>C</sub> and E[cost]<sub>I</sub> for 100 instances)



#### 4. Nurse staffing decisions under various pool structures.

• Result



One pool: lowest staffing cost (21 cross-training)
 Optimal disjoint pool : lowest staffing cost (2 cross-training)

- 5. Which units to pool together?
  - Motivation
    - In hospital, each unit has different level of uncertainties



- Which units to pool together?

- 5. Which units to pool together?
  - Experiment
    - Label units with *low uncertainties* as "A".
    - Label units with *high uncertainties* as "B".
    - Consider 3 cases:
      - (case 1) demand only.
      - (case 2) absenteeism only.
      - (case 3) both demand and absenteeism.
    - For each case, we solve the ONPD model.
    - Resulting pools: (Type 1) *A units only*. (Type 2) *B units only*. (Type 3) *both A and B units*.



- 5. Which units to pool together?
  - Result (Expected number of Type 1, 2, 3 pools)



Units with high uncertainties are grouped together (Type 2 pool).
Absenteeism plays an important role on the pooling strategy.

# Thank You for Your Attention! Questions / Comments?