

Shift Scheduling in Pediatric Emergency Medicine

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Content

- Background
- Motivation
- Formulations
 - Weighted Sum method
 - Metric constraints method
- Result
- Future Research
 - Pareto method

Resident Responsibilities in the U-M Pediatric Emergency Department

- 3-7 year medical training program
 - Responsibilities differ by residency year
- Balancing patient care and educational requirements
 - In hospital
 - Caring for patients
 - Teaching medical students
 - Learning from attending physicians
 - Out of hospital
 - Community clinics
 - Conferences
 - Other educational requirements

Pediatric ED: Scheduling Considerations

- All shifts assigned to a resident
- Appropriate coverage
 - e.g. certain shifts require a senior resident
- ACGME rules (similar to ABET for engineering)
 - e.g. 10 hour break rule
- Several different residency programs
 - Pediatrics (PED)
 - Family practice (FP)
 - Emergency medicine (EM)
- And others

Motivation

- **Scheduling residents**
 - **Complicated requirements (UM Pediatric ED)**
 - 25 governing rules and preferences
 - Educational goals
 - Patient care
 - Regularization / Safety
 - **Chief resident built monthly schedule by hand**
 - Time consuming process: 20 - 25 hours / month
 - Transfer every year: no scheduling experience in July
 - Guess and check: errors / tedious correction process

	3			1		7	
6			8				2
		1		4		5	
	7				2		4
2				9			6
	4		3				1
		5		3		4	
1					6		5
	2		1			3	



Mixed Integer Programming

Motivation

- **Practical Significance**

- Poor-quality schedule

- Residents: decreased interest in learning
 - Patients: adverse health events

(Smith-Coggins R, et. al. (1994) : "Relationship of day versus night sleep to physician performance and mood." Annals of Emergency Medicine)

- **Goals**

- Solves for feasible schedule quickly
 - Create a good quality schedule with no violations



Formulation: Problem Size

- Sets**

- R: set of residents
 - 15-25 residents
- D: set of days in the schedule
 - 35 days
- S: set of shifts
 - 8 shifts

- Decision Variables**





- Binary: $x_{rds} \in \{0, 1\}$
 - 1 if resident r works shift s on day d
 - 0 otherwise

Residents Name					
Smith	Sanchez	Chen	Shah	...	

	27 th	...	1 st	...	31 st
7a-4p	Shah	
9a-6p	Joe	Shah
10a-7p		
12p-9p	Chen	Chen
4p-1a	Smith	...	Sanchez	...	
5p-2a		Sanchez
8p-5a	Sanchez	...	Smith	...	Smith
11p-8a		...	Chen	...	Joe

Objectives: Shift Fairness

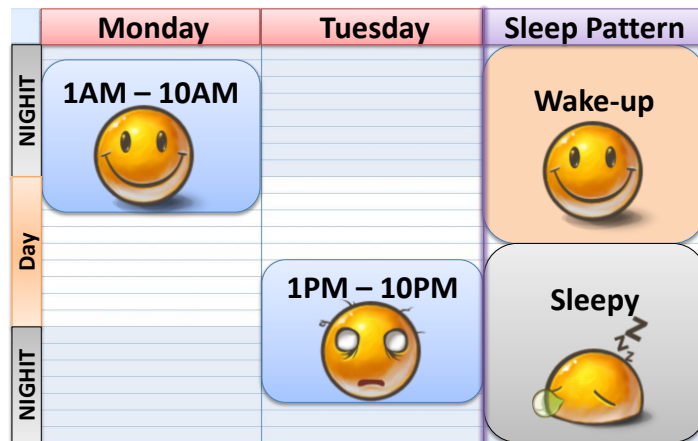
- **Total / night shift equity**
 - Equal opportunities for training
 - Improved morale and learning ability

Resident Name	Smith	Jones	Chen	Joe
Night Shifts / Total Shifts	0 / 7	1 / 7	1 / 7	5 / 7
Fairness				

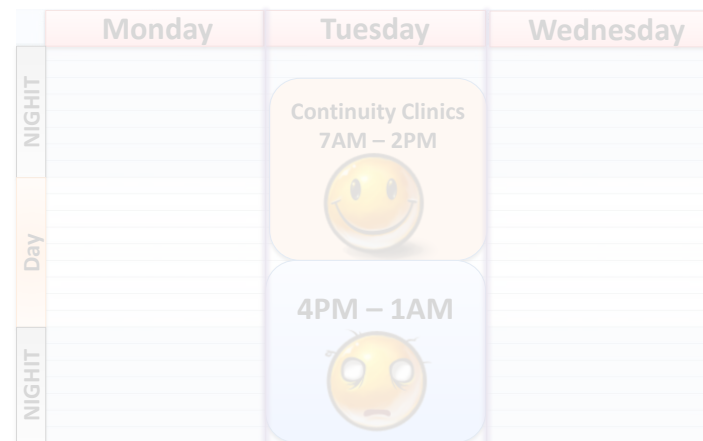
- **Total shift equity (TSE):** $(\sum t_{ij}, t_{ij} = |D_i - D_j|, i > j)$
- **Night shift equity (NSE):** $(\sum n_{ij}, n_{ij} = |N_i - N_j|, i > j)$

Objectives: Undesired Shift

- **Bad sleep patterns and post-clinic shifts**
 - Improves resident quality of life
 - Increases patient safety



Bad sleep pattern







Post-Clinic shift



- **Minimum bad sleep patterns (BSP):** $(\sum count)$

Objectives: Undesired Shift

- **Bad sleep patterns and post-clinic shifts**
 - Improves resident quality of life
 - Increases patient safety

	Monday	Tuesday	Sleep Pattern
NIGHT	1AM – 10AM 		Wake-up 
Day			
NIGHT		1PM – 10PM 	Sleepy 

Bad sleep pattern

	Monday	Tuesday	Wednesday
NIGHT			
Day		Continuity Clinics 7AM – 2PM 	
NIGHT		4PM – 1AM 	

Post-Clinic shift

- **Minimum post-clinic shifts (PCC):** $(\sum count)$

Formulation: Constraints

- **Constraints (rules/requirements)**
 - One resident assigned to each shift in the month
 - $\sum_{r \in \{\text{all}\}} x_{rds} = 1, \quad \forall d, \forall s$
 - Meets shift requests
 - $x_{rds} = 0, \quad \forall r, \forall d, s \in \{\text{day off, conferences, continuity clinic}\}$
 - Ensure resident type appropriate for shift
 - $\sum_{r \in \{\text{PED}\}} \sum_{s \in P} x_{rsd} \geq 1, \forall d, P = \{\{7a, 9a\}, \{4p, 5p\}, \{8p, 11p\}\}$
 - Intern-forbidden shifts
 - $\sum_{r \in \{\text{interns}\}} \sum_d x_{rsd} = 0, \forall s \in \{7a, 11p\}$
 - And others

Multi-Criteria Problem

- **Multi-Criteria schedule**
 - **Metrics for UM Pediatric Emergency Department**
 - Total shift equity (TSE)
 - Night shift equity (NSE)
 - Minimum bad sleep patterns (BSP)
 - Minimum post-clinic shifts (PCC)
 - \vdots
- 

Multi-objective Mathematical Programming

Weighted Sum Method

$$\begin{aligned} \text{Min } & \mathbf{w_1}(TSE) + \mathbf{w_2}(NSE) + \mathbf{w_3}(BSP) + \mathbf{w_4}(PCC) \\ \text{s. t. } & \text{"rules/requirements"} \\ & x_{rds} \in \{0,1\} \end{aligned}$$

- Quantifying preferences (w_i) is difficult
 - Weights are subjective and difficult to quantify
 - Resulting schedule does not match their intentions
 - Various measurement units
 - Equity (σ , $\text{Max}|\text{diff}_{ij}|$, $\sum|\text{diff}_{ij}|$, ...)

Optimized Residency Scheduling Assistant (ORSA): Metrics Formulation

- Feasibility problem
 - Constraint on metrics

$$\begin{array}{ll}\text{Min } w_1(TSE) + w_2(NSE) + w_3(BSP) + w_4(PCC) \\ \text{s. t.} & \text{"rules/requirements"} \\ & x_{rds} \in \{0,1\}\end{array}$$

- Benefits of a feasibility problem
 - More flexible
 - Faster to solve: < 2 sec.
 - Given: 35 days / 20 PEDs / 8 shifts

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Optimized Residency Scheduling Assistant (ORSA) : Interactive Improvement

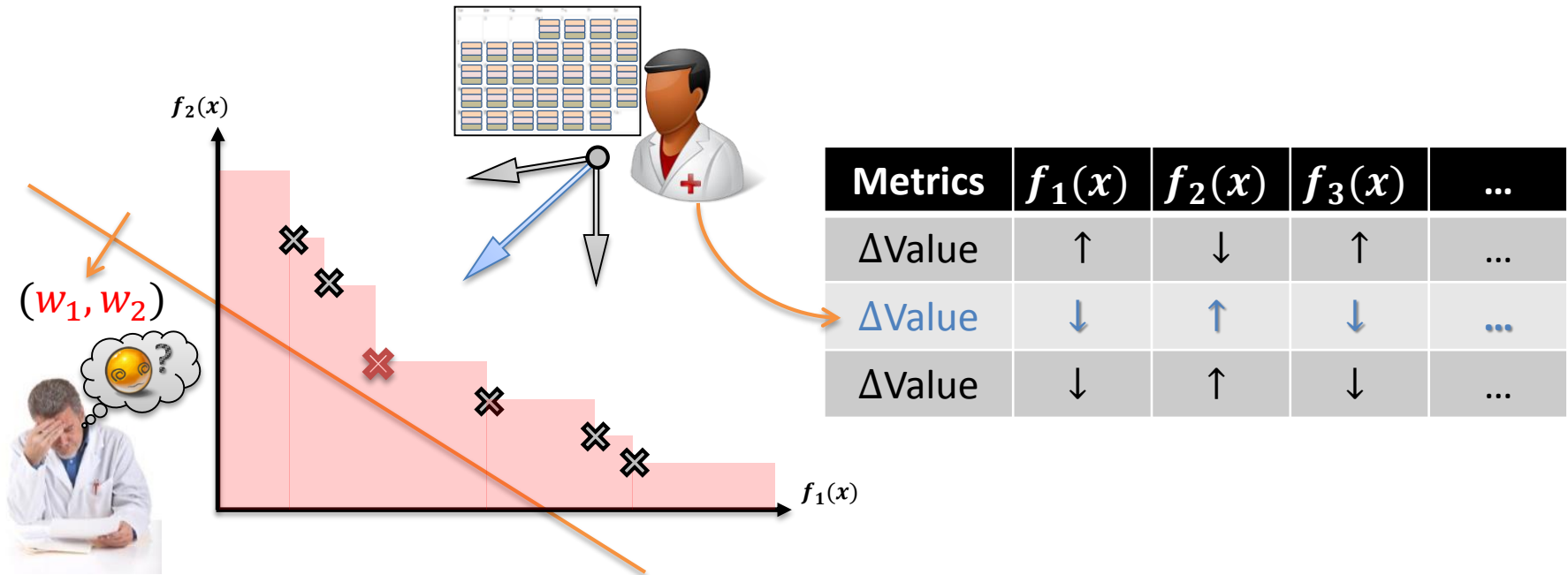
- Example output of metrics
 - Value (Lower bound, Upper bound)

Resident Name	Number of Shifts	Number of Night Shifts	Number of Post CC	Number of Bad Sleep Templates
Smith	8 (7,9)	2 (0,10)	0 (0,1)	1 (0,1)
Sanchez	8 (7,10)	1 (0,10)	0 (0,1)	1 (0,1)
Chen	8 (7,9)	5 (0,10)	1 (0,1)	1 (0,1)
Shah	14 (13,15)	3 (0,10)	1 (0,1)	1 (0,1)
:	:	:	:	:

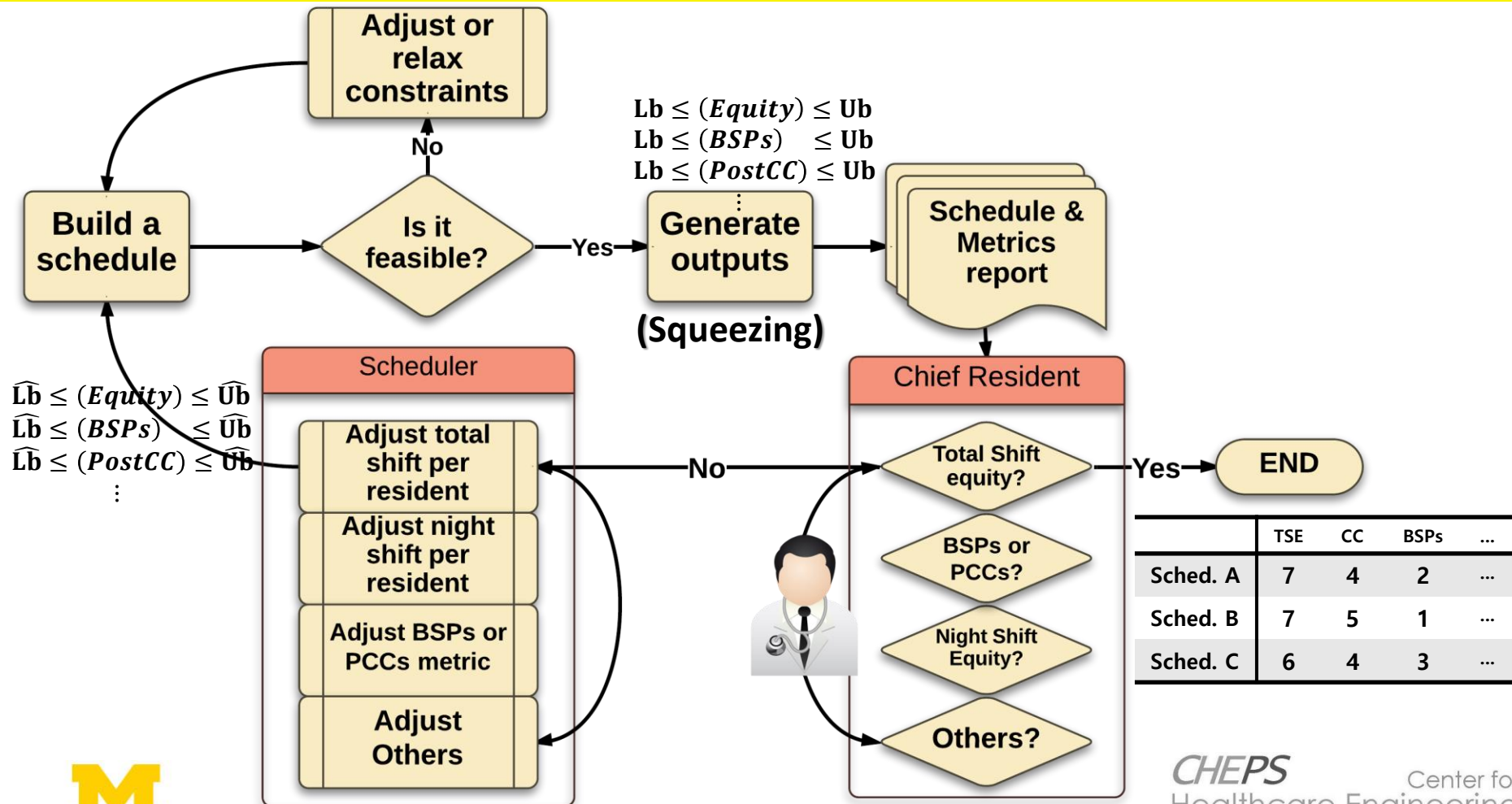
- Interactive approach engaging chief resident
 - Iteratively adjust bounds on metric constraints
 - Quickly build high quality schedule

Optimized Residency Scheduling Assistant (ORSA) : Interactive Improvement

- Interactive Feedback
 - Chief resident identifies undesirable qualitative characteristics

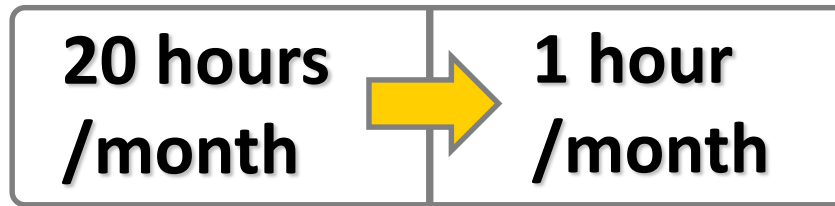


ORSA Methodology

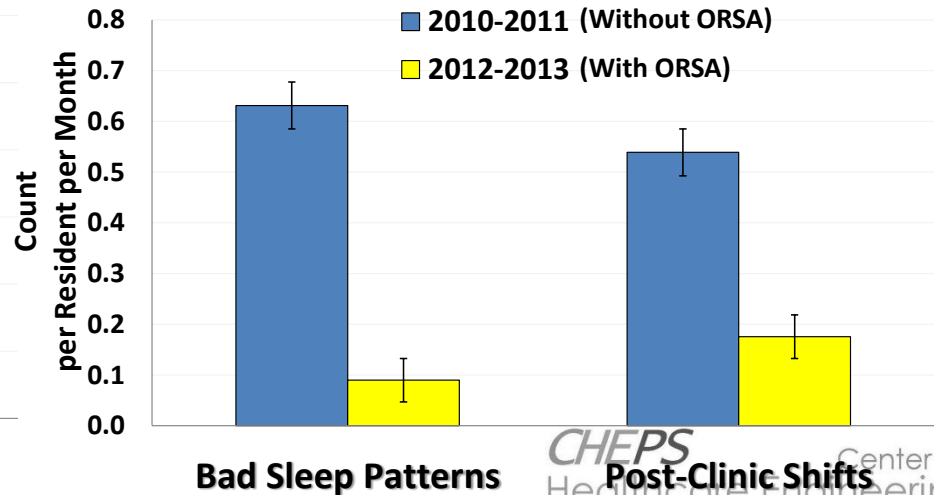
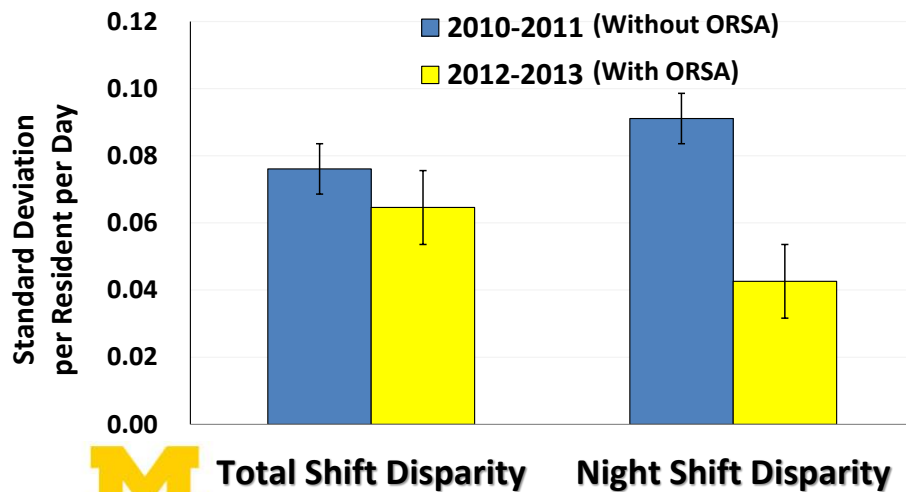


ORSA: Results

- Our metrics-based scheduling tool:
 - Reduces time to create schedules



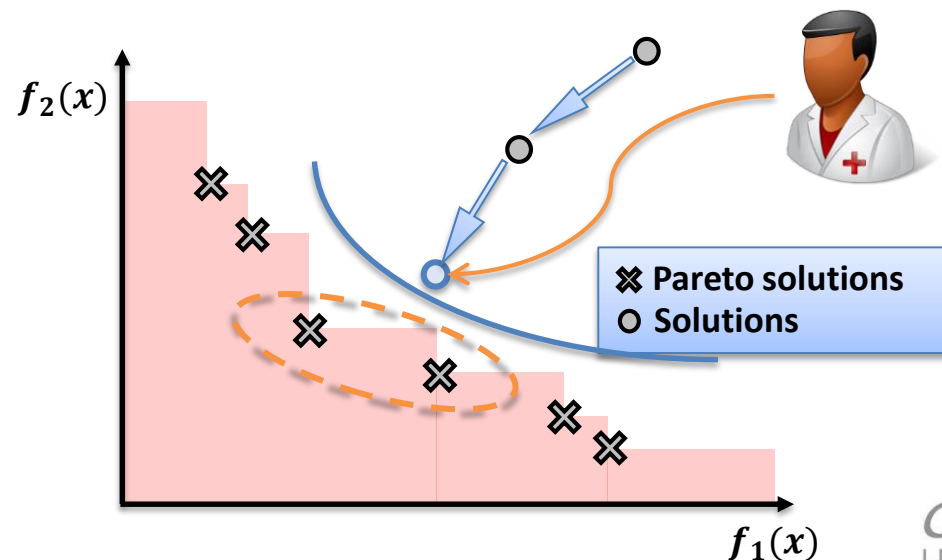
- Solves a multi-criteria scheduling problem



ORSA: Limitations

- **Myopic Solution**

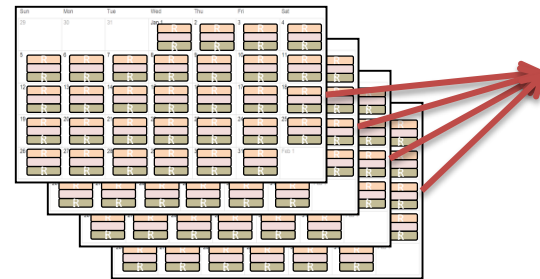
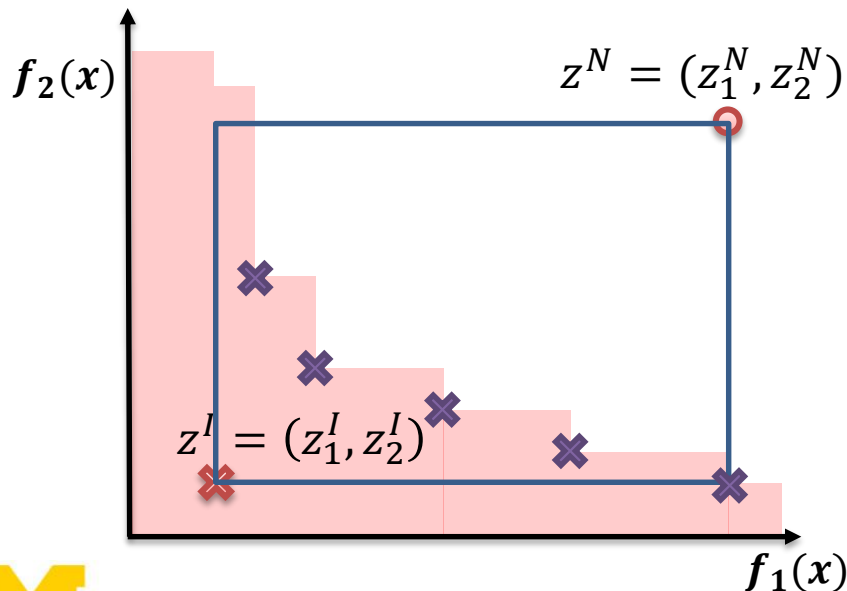
- Non-Pareto solution could be selected by a chief residents
 - Never see the whole picture (the set of Pareto solutions)
 - The most preferred solution is “most preferred” with respect to their satisfaction



Next Step

- **Pareto Solutions**

- Generate the Pareto solutions of the problem (all of them or a sufficient representation)
 - Select the most preferred one among them



Efficient Schedules



Pareto: Bi-Objective Problem

- **Notation**

- \mathcal{H} : Solution Space, the set of feasible solutions
- \mathcal{P} : Pareto Set
- $z_i = f_i(x)$: i th integer objective function, $\in \mathbb{Z}$
- Dominance ($<$) : $x < x'$ if and only if $z_i \leq z_i'$ where at least one inequality is strict

- **Bi-Objective Problem**

$$\begin{aligned} \min f(x) &= (f_1(x), f_2(x)) \\ \text{s.t. } x &\in \mathcal{H} \end{aligned}$$

Pareto: Bi-Objective Problem

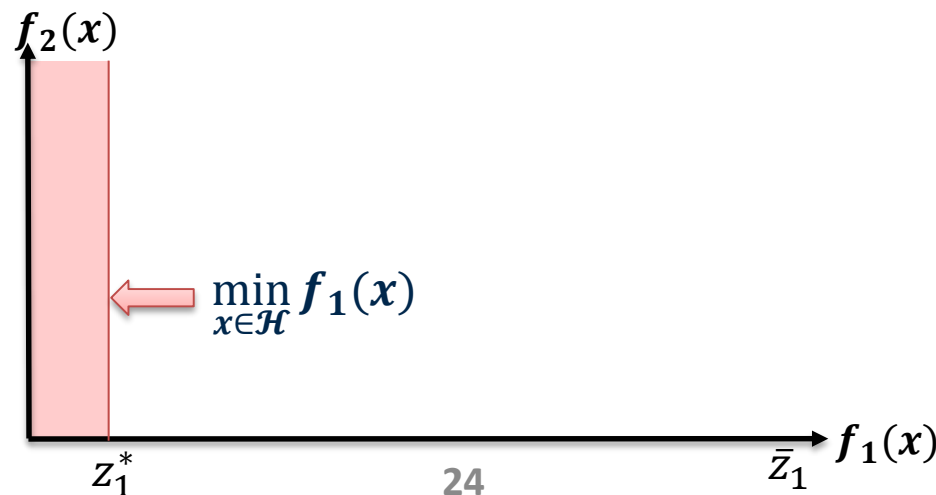
- Pareto Square Region

- Ideal Point:

- $z_1^* = \min_{x \in \mathcal{H}} f_1(x)$ and $z_2^* = \min_{x \in \mathcal{H}} f_2(x)$

- Nadir Point:

- $\bar{z}_1 = \min_{x \in \mathcal{H} \cap f_2(x)=z_2^*} f_1(x)$ and $\bar{z}_2 = \min_{x \in \mathcal{H} \cap f_1(x)=z_1^*} f_2(x)$



Pareto: Bi-Objective Problem

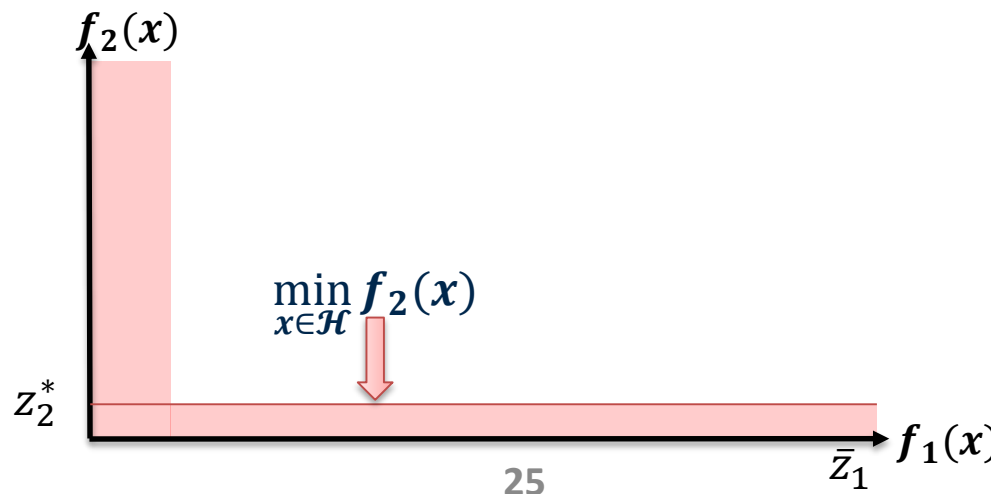
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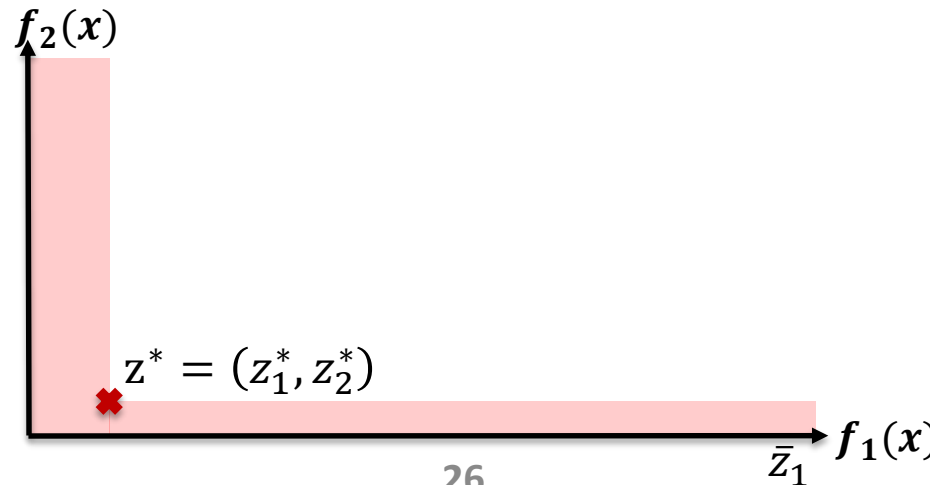
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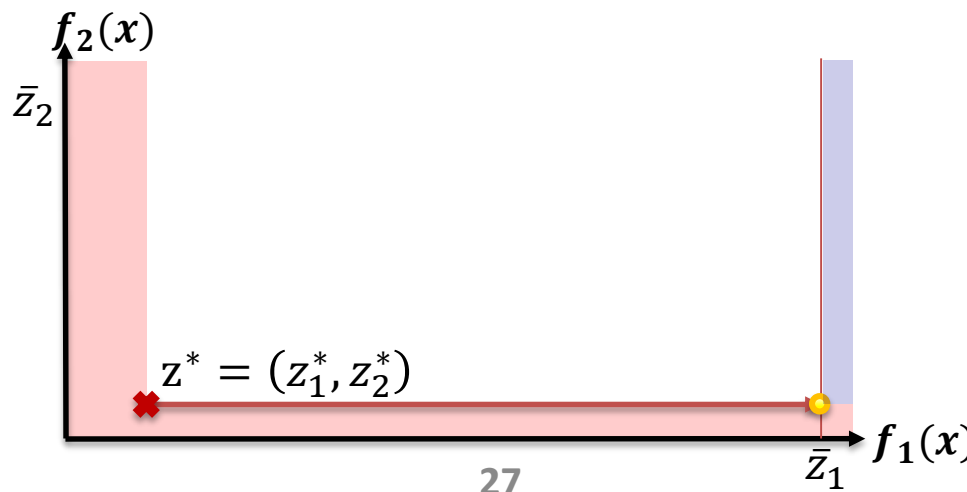
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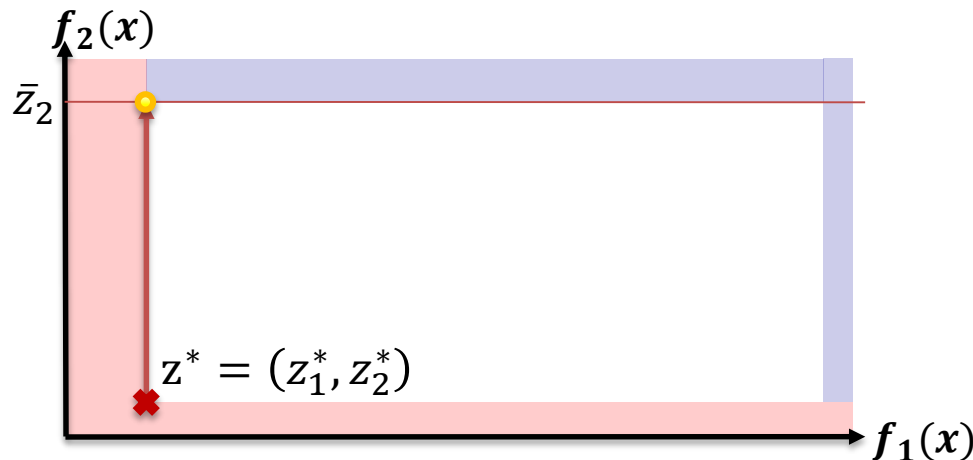
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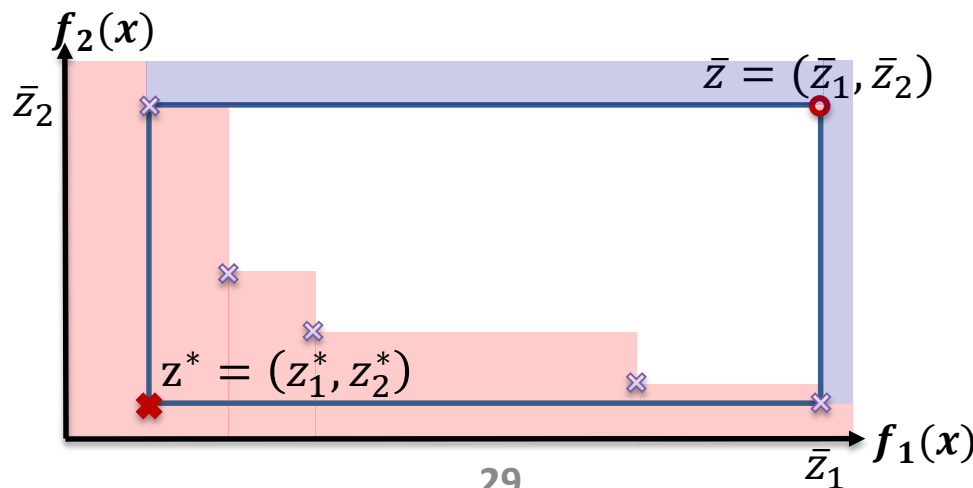
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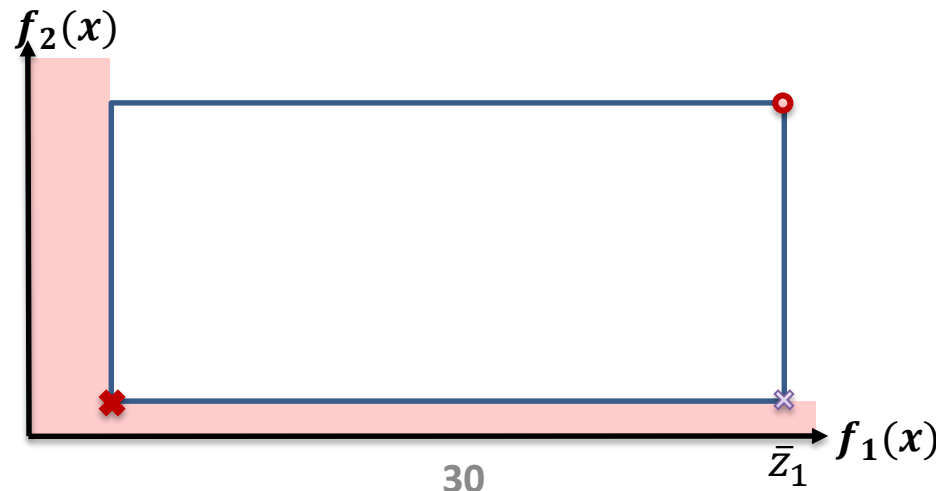


Pareto: Bi-Objective Problem

- Pareto Squeezing Algorithm

Algorithm 1 Exact squeezing algorithm for bi-objective problems

```
Let  $P$  is set of pareto solutions we've found;  
Compute the ideal  $(z_1^*, z_2^*)$  and Nadir  $(\bar{z}_1, \bar{z}_2)$  points;  
Set  $P := \{(\bar{z}_1, z_2^*)\}$  and  $\delta := \bar{z}_1 - 1$ ;  
WHILE  $\delta \geq z_1^*$   
    Solve  $P_2(\delta)$  and get optimal solution  $(z_1^i, z_2^i)$  to  $P_2(\delta)$ ;  
    //Given  $(z_1^i, z_2^i)$ , Find a left-bottom corner  $(\hat{z}_1^i)$  in the Pareto set;  
    Solve  $SQZ_1(z_2^i)$  and get optimal solution  $(\hat{z}_1^i, z_2^i)$  to  $SQZ_1(z_2^i)$ ;  
END WHILE  
Set  $P := P + (\hat{z}_1^i, z_2^i)$  and  $\delta = \hat{z}_1^i - 1$ ;  
GO Step 2 UNTIL  $z_1 = z_1^*$ ;
```



Pareto: Bi-Objective Problem

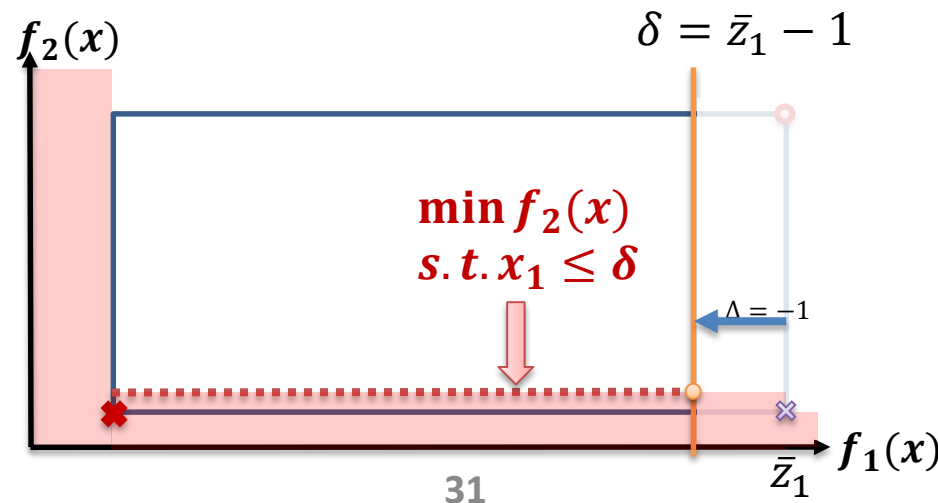
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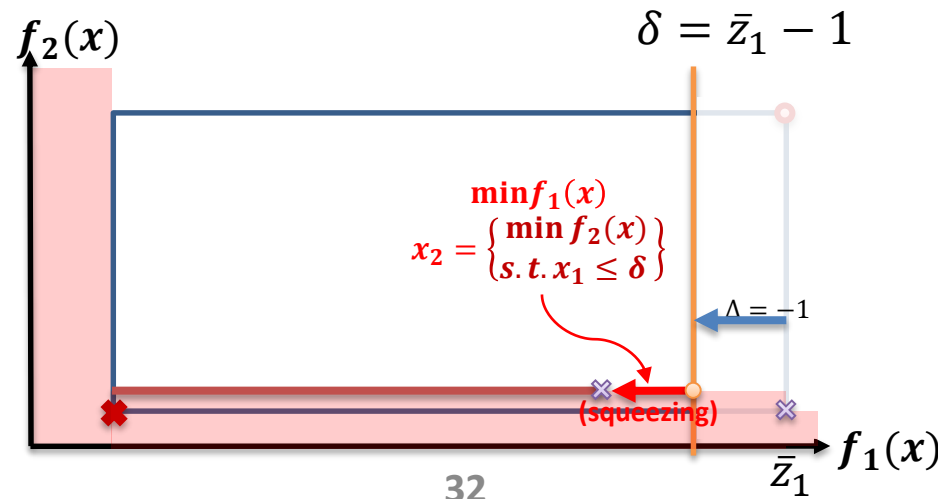
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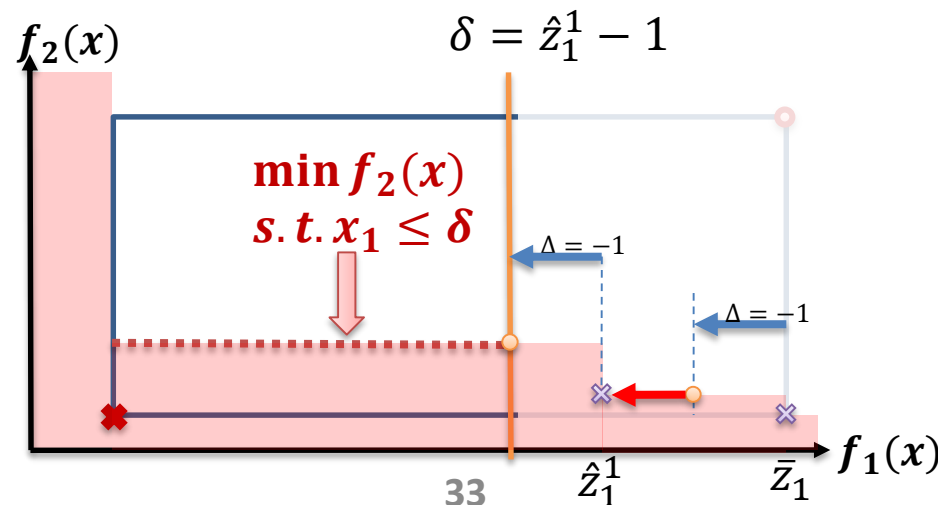
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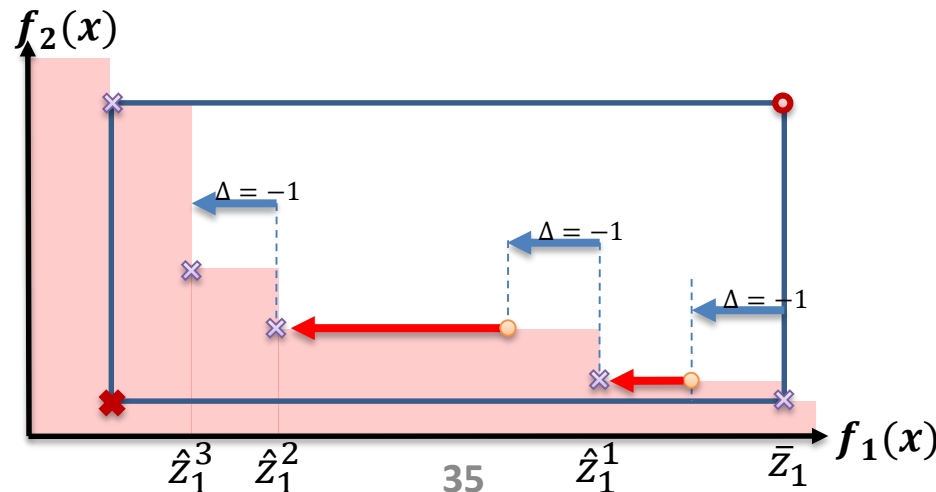
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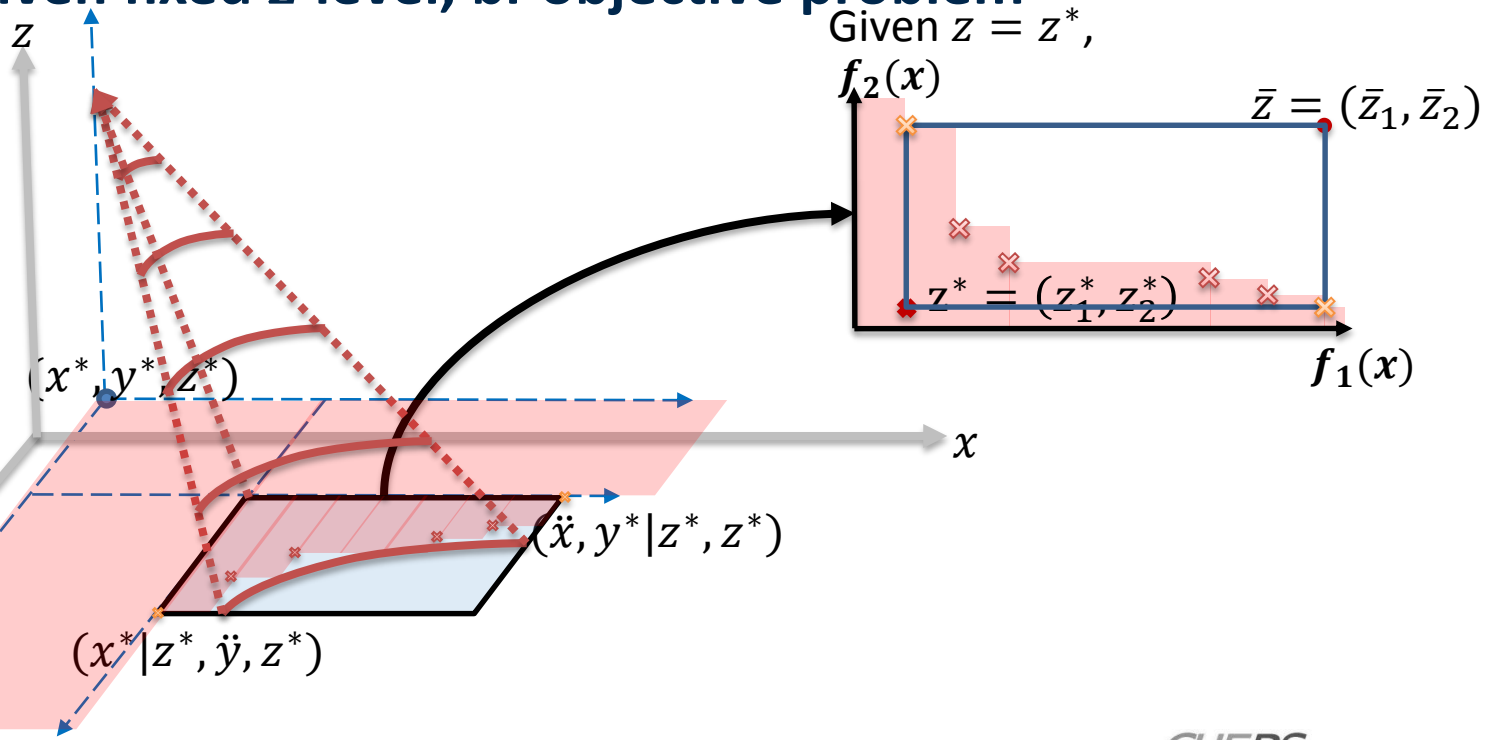
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Pareto: Tri-Objective Problem

- “Pareto Cone”

- Given fixed z level, bi-objective problem



Acknowledgements

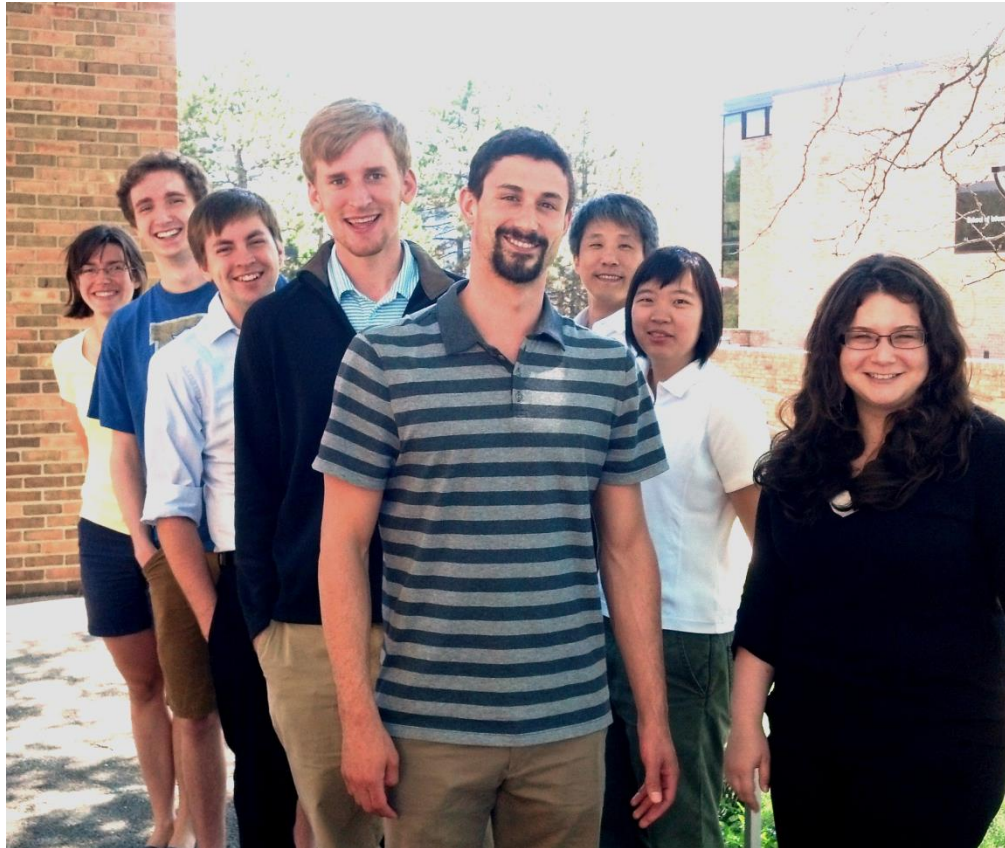
- Thank you to CHEPS, TDC Foundation, the Bonder Foundation, and Dr. Brian Jordan, Dr. Micah Long, Dr. Jenny Zank and Dr. Ed O'Brien for making this research possible.

CHEPS and the HEPS Master's Program

- **CHEPS:** The Center for Healthcare Engineering and Patient Safety
- **HEPS:** Industrial and Operations Engineering (IOE) Master's Concentration in Healthcare Engineering and Patient Safety offered by CHEPS
- CHEPS and HEPS offer unique multidisciplinary teams from engineering, medicine, public health, nursing, and more collaborating with healthcare professionals to better provide and care for patients
- For more information, contact Amy Cohn at amycohn@umich.edu or visit the CHEPS website at: <https://www.cheeps.engin.umich.edu>



Thank You!



Feedback and Questions

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University of Michigan*