

# Block Scheduling for a Pediatric Residency Program

Peter Mayoros, B.S.E.

University of Michigan

October 3, 2014



# My Collaborators

- Research Team:
  - Amy Cohn, Ph.D.
  - Zak Verschure
  - Young-Chae Hong
  - Ji Wang
  - William Pozehl
- Contacts at the UM Health System:
  - Edward O’Brien, M.D.
  - Jenny Shin, M.D.

# Presentation Outline

- Background
- Motivation
- Model Formulation
- Model Implementation
- Results
- Future Work



# Pediatrics at Michigan

- Based in C.S. Mott Children's Hospital
- Around 100 residents at Motts every year
- The hospital is very diverse in the programs it offers



# Services in Pediatrics

- There are many different areas in Mott that residents must gain experience in:
  - E.g. Neonatal Intensive Care Unit (NICU), General Care, Hematology & Oncology (HemOnc)
- Every month each of these services must be adequately staffed and there is a minimum and maximum number of residents allowed in each service



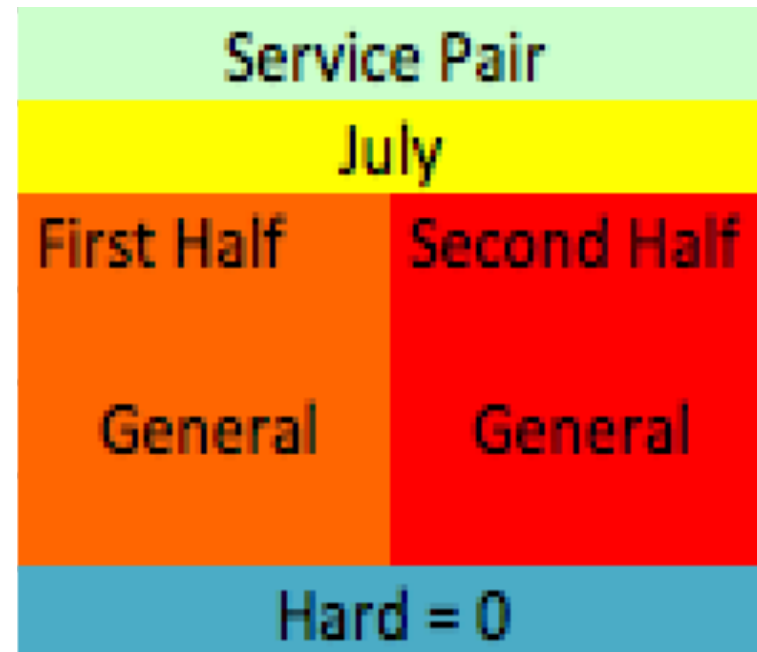
# Residency at Michigan

- After four years in medical school the students are “matched” to a residency program
- These first year residents are called “interns”
- After the first year they become “senior residents”
- Every resident has a minimum and maximum amount of time (in half months) that they must spend in each service
- We were tasked with compiling a block schedule for just the interns



# Service Pairs

- A couplet of services that can be worked during the same month.
- The first service in the pair is worked for the first half of the month and the second service is worked for the second half of the month
- Different examples of service pairs could be:
  - General/general or NICU/general
- Different combinations can also be determined as hard or not



# What is a Block Schedule?

Joe Smith	Service Pair 1		Service Pair 2		Service Pair 3	
	July		August		September	
	First Half	Second Half	First Half	Second Half	First Half	Second Half
	General	General	NICU	HemOnc	General	Vacation
Hard = 0		Hard = 1		Hard = 0		
Jill O'Brien	Service Pair 4		Service Pair 1		Service Pair 2	
	July		August		September	
	First Half	Second Half	First Half	Second Half	First Half	Second Half
	HemOnc	NICU	General	General	NICU	HemOnc
Hard = 1		Hard = 0		Hard = 0		

Assigning service pairs to residents over the course of a year

NOTE: vacation as a service



# Motivation

- Before:
  - Block Schedule was compiled by hand by the chief resident
    - This took a lot of time
    - If changes needed to happen after a draft was completed they had to rewrite the entire schedule
- Now:
  - A computer program compiles the schedule in less than 5 minutes
  - Changes are effortless



# Model Overview

## Sets:

$R$  : set of all residents that need to be scheduled

$P$  : set of all service pairs

$S$  : set of all services that need residents

$M$  : set of months

$D$  : set of desires. this is a list given to us by the chief resident that details the residents vacation requests

$p_{si} \subset P$  subset of service pairs  $P$  that have service  $S$  in half-month  $i$  where  $i \in \{0, 1\}$

$h_p \subset P$  : set of pre-defined hard service pairs

$C_d \subset D$  : set of candidate assignments for desire  $D$



# Model Overview

## Parameters:

$q_{ps}$  : the number of half-months of service  $s$  that are in service pair  $p$   $\forall s \in S, p \in P$

$b_p$  : if service pair  $p$  is difficult 1, if easy 0,  $\forall p \in P$

$u_{sm}$  : an upper bound for the number of residents service  $s$  needs in month  $m$   $\forall s \in S, m \in M$

$l_{sm}$  : a lower bound for the number of residents service  $s$  needs in month  $m$   $\forall s \in S, m \in M$

$u_{rs}$  : a upper bound for the number of months required on service  $s$  for the year  $\forall s \in S, r \in R$

$l_{rs}$  : a lower bound for the number of months required on service  $s$  for the year  $\forall s \in S, m \in M$

$l_d$  : a lower bound for the number of desires filled for month  $\forall m \in M$

$u_d$  : an upper bound for the number of desires filled for month  $\forall m \in M$



# Model Overview

## Decision Variables:

$X_{rpm}$  : Whether or not resident  $r$  is scheduled to service pair  $p$  on month  $m \quad \forall r \in R, p \in P, m \in M$

$Y_{rm}$  : Whether or not resident  $r$  works 3 or more hard services in a row beginning on month  $m \quad \forall r \in R$

$Z$  : Helper variable for maximum number of times a resident works 3 difficult shifts in a row.

## Objective Function:

Constant

## Variable Restrictions:

$$\begin{aligned} X_{rtm} &\in \{0,1\} \\ Y_{rm} &\in \{0,1\} \\ Z &\in (R) \end{aligned}$$



# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

Each Resident is assigned one service pair per month



# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

$$l_{sm} \leq \sum_{p \in P_{si}} X_{rpm} \leq u_{sm} \quad \forall m \in M, s \in S, i \in 0, 1$$

Each service has at least the minimum number of residents but no more than the Maximum, for all half months



# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

$$l_{sm} \leq \sum_{p \in P_{si}} X_{rpm} \leq u_{sm} \quad \forall m \in M, s \in S, i \in 0, 1$$

$$l_{rs} \leq \sum_{p \in P} \sum_{m \in M} q_{ps} * X_{rpm} \leq u_{rs} \quad \forall r \in R, s \in S$$

Each resident has at least the minimum number of months in each service but no more than the maximum, for all months



# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

$$l_{sm} \leq \sum_{p \in P_{si}} X_{rpm} \leq u_{sm} \quad \forall m \in M, s \in S, i \in 0, 1$$

$$l_{rs} \leq \sum_{p \in P} \sum_{m \in M} q_{ps} * X_{rpm} \leq u_{rs} \quad \forall r \in R, s \in S$$

$$b_t X_{rpm} + b_t X_{rp(m+1)} + b_t X_{rp(m+2)} \leq Y_{rm} + 2$$

Y is equal to 1 if resident r works more than 3 'hard' service pairs in a row





# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

$$l_{sm} \leq \sum_{p \in P_{si}} X_{rpm} \leq u_{sm} \quad \forall m \in M, s \in S, i \in 0, 1$$

$$l_{rs} \leq \sum_{p \in P} \sum_{m \in M} q_{ps} * X_{rpm} \leq u_{rs} \quad \forall r \in R, s \in S$$

$$b_t X_{rpm} + b_t X_{rp(m+1)} + b_t X_{rp(m+2)} \leq Y_{rm} + 2$$

$$Z \geq \sum_{m \in M} Y_{rm} \quad \forall r \in R$$

Z is equal to the maximum number of times a resident works 3 difficult services in a row



# Model Overview

$$\sum_{p \in P} X_{rpm} = 1 \quad \forall r \in R, m \in M :$$

$$l_{sm} \leq \sum_{p \in P_{si}} X_{rpm} \leq u_{sm} \quad \forall m \in M, s \in S, i \in 0, 1$$

$$l_{rs} \leq \sum_{p \in P} \sum_{m \in M} q_{ps} * X_{rpm} \leq u_{rs} \quad \forall r \in R, s \in S$$

$$b_t X_{rpm} + b_t X_{rp(m+1)} + b_t X_{rp(m+2)} \leq Y_{rm} + 2$$

$$Z \geq \sum_{m \in M} Y_{rm} \quad \forall r \in R$$

$$l_d \leq \sum_{c \in C_d} X_{r(c)p(c)m(c)} \leq u_d \quad \forall d \in D$$

The amount of desired vacation requests filled in a schedule must be in between the upper and lower bounds

# Model Implementation

- We coded the model into a c++ program with the help of the cplex library
- The Chief resident submitted input files specifying the interns, services, service pairs and their upper and lower bounds
- Our program read in the input files and then using their data produced a block schedule for the pediatric department interns



# Results

- We created a block schedule for 27 interns for this current year in under 5 minutes per run
- We saved the chief resident countless hours in schedule making
- We found a more optimal schedule than the chief resident could have on her own



# Future Work

- We want to integrate all types of resident regardless of their year in the program so that we can optimize the total schedule
- We are restructuring our model to make it more user friendly and easier to understand



# Questions?

Peter Mayoros

[pmayoros@umich.edu](mailto:pmayoros@umich.edu)

Prof. Amy Cohn

[amycohn@med.umich.edu](mailto:amycohn@med.umich.edu)

*Department of Industrial and Operations Engineering*

*University of Michigan*

