

Pareto Optimality in Pediatric Residency Shift Scheduling

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University of Michigan

ISERC at Nashville (05 / 31 / 2015)

Collaborators

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- **Edmond O'Brien, M.D**
- **William Pozehl, M.S.E.**
- **Zachary VerSchure**

Content

- **Background**
- **Motivation**
- **Formulations**
 - **Weighted sum method**
 - **Metric constraints method**
- **Result**
- **Ongoing Research**
 - **Pareto method**



Resident Responsibilities in the U-M Pediatric Emergency Department

- **3-7 year medical training program**
 - Responsibilities differ by residency year
- **Balancing patient care and educational requirements**
 - **In hospital**
 - Caring for patients
 - Teaching medical students
 - Learning from attending physicians
 - **Out of hospital**
 - Community clinics
 - Conferences
 - Other educational requirements



Pediatric ED: Scheduling Considerations

- **All shifts assigned to a resident**
- **Appropriate coverage**
 - e.g. certain shifts require a senior resident
- **ACGME rules (similar to ABET for engineering)**
 - e.g. 10 hour break rule
- **Several different residency programs**
 - Pediatrics (PED)
 - Family practice (FP)
 - Emergency medicine (EM)
- **And others**



Motivation

- **Scheduling residents**

- **Complicated requirements**

- **25 governing rules and preferences**

- Educational goals

- Patient care

- Regulations / Safety

	3			1		7
6			8			2
	1		4		5	
	7			2		4
2			9			6
	4		3			1
		5		3		4
1				6		5
	2		1			3



- **Chief resident built monthly schedule by hand**

- Time consuming process: 20 - 25 hours / month

- Transfer every year: no scheduling experience in July

- Guess and check: errors / tedious correction process

Motivation

- **Practical Significance**

- Poor-quality schedule

- **Residents: decreased interest in learning**

- **Patients: adverse health events**

(Smith-Coggins R, et. al. (1994) : "Relationship of day versus night sleep to physician performance and mood." Annals of Emergency Medicine)

- **Goals**





- Solve for feasible schedule quickly

- Create a good quality schedule with no violations



Objectives: Shift Fairness

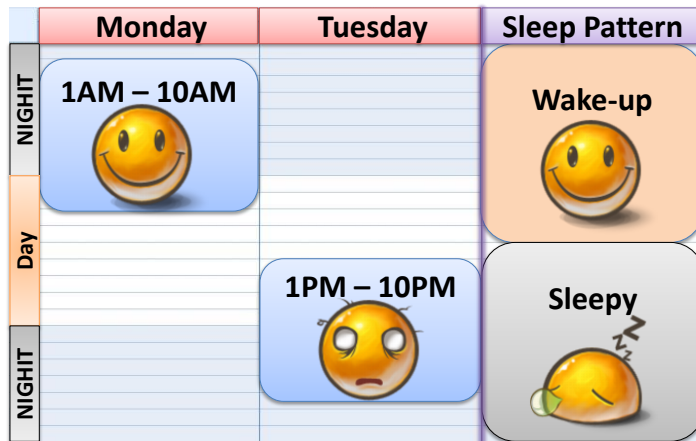
- **Total / night shift equity**
 - Equal opportunities for training
 - Improved morale and learning ability

Resident Name	Smith	Jones	Chen	Joe
Night Shifts / Total Shifts	0 / 7	1 / 7	1 / 7	5 / 7
Fairness				

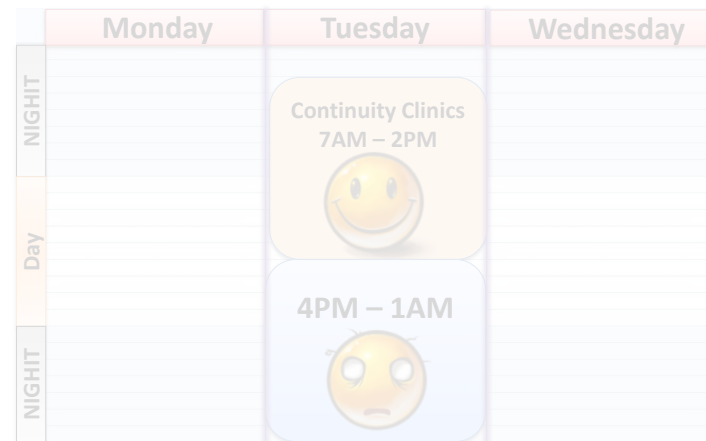
- **Total shift equity (TSE):** $(\sum t_{ij}, t_{ij} = |D_i - D_j|, i > j)$
- **Night shift equity (NSE):** $(\sum n_{ij}, n_{ij} = |N_i - N_j|, i > j)$

Objectives: Undesired Shift

- **Bad sleep patterns and post-clinic shifts**
 - Reduces resident quality of life
 - Decreases patient safety



Bad sleep pattern

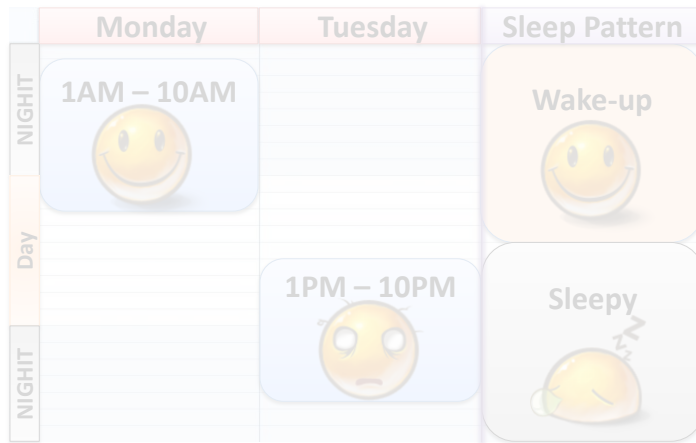


Post-Clinic shift

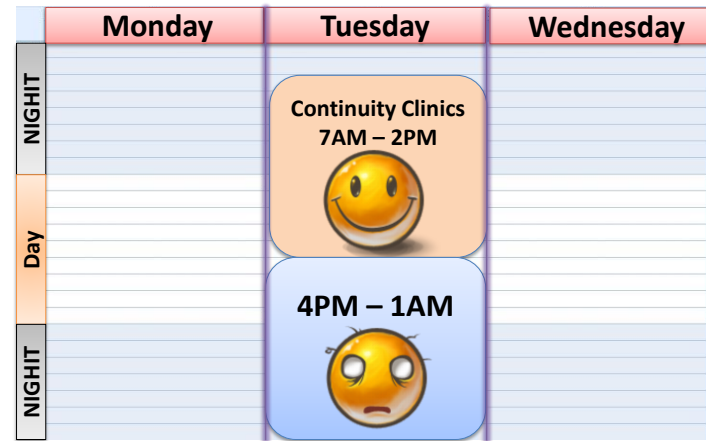
- **Minimum bad sleep patterns (BSP):** $(\sum count)$

Objectives: Undesired Shift

- **Bad sleep patterns and post-clinic shifts**
 - Reduces resident quality of life
 - **Decreases patient safety**



Bad sleep pattern



Post-Clinic shift

- **Minimum post-clinic shifts (PCC):** $(\sum count)$

Formulation: Decision Variables

- **Sets**

- **R**: set of residents
- **D**: set of days in the schedule
- **S**: set of shifts

- **Decision Variables**

- **Binary**: $x_{rds} \in \{0, 1\}$
 - 1 if resident r works shift s on day d
 - 0 otherwise

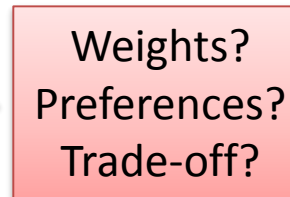
Residents Name					
Smith	Sanchez	Chen	Shah	...	
	27 th	...	1 st	...	31 st
7a-4p	Shah	
9a-6p	Joe	Shah
12p-9p	Chen	Chen
4p-1a	Smith	...	Sanchez	...	
5p-2a		Sanchez
8p-5a	Sanchez	...	Smith	...	Smith
11p-8a		...	Chen	...	Joe

Formulation: Constraints

- **Constraints (rules/requirements)**
 - **One resident assigned to each shift in the month**
 - $\sum_{r \in \{\text{all}\}} x_{rds} = 1, \forall d, \forall s$
 - **Meets shift requests**
 - $x_{rds} = 0, \forall r, \forall d, s \in \{\text{day off, conferences, continuity clinic}\}$
 - **Ensure resident type appropriate for shift**
 - $\sum_{r \in \{\text{PED}\}} \sum_{s \in P} x_{rsd} \geq 1, \forall d, P = \{\{7a, 9a\}, \{4p, 5p\}, \{8p, 11p\}\}$
 - **Intern-forbidden shifts**
 - $\sum_{r \in \{\text{interns}\}} \sum_d x_{rsd} = 0, \forall s \in \{7a, 11p\}$
 - **And others**

Multi-Criteria Problem

- **Multi-Criteria schedule**
 - **Metrics for UM Pediatric Emergency Department**
 - Total shift equity (TSE)
 - Night shift equity (NSE)
 - Minimum bad sleep patterns (BSP)
 - Minimum post-clinic shifts (PCC)
 - \vdots



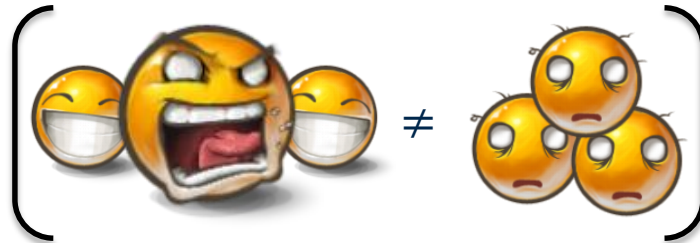
Weights?
Preferences?
Trade-off?

Multi-objective Mathematical Programming

Weighted Sum Method

$$\begin{aligned} \text{Min } & w_1(TSE) + w_2(NSE) + w_3(BSP) + w_4(PCC) \\ \text{s. t. } & \text{"rules/requirements"} \\ & x_{rds} \in \{0,1\} \end{aligned}$$

- **Quantifying preferences (w_i) is difficult**
 - **Weights are subjective and difficult to quantify**
 - Resulting schedule does not match their intentions
 - **Various measurement units**
 - **Equity** (σ , $\text{Max}|\text{diff}_{ij}|$, $\sum|\text{diff}_{ij}|$, ...)
 - **Non-linearity**
 - **10 BSPs \neq 10 x 1 BSP**



Metrics Formulation

- **Feasibility problem**
 - **Constraint on metrics**

$$\begin{array}{ll} \text{Min } w_1(TSE) + w_2(NSE) + w_3(BSP) + w_4(PCC) \\ \text{s. t.} & \text{"rules/requirements"} \\ & x_{rds} \in \{0,1\} \end{array}$$

- **Benefits of a feasibility problem**
 - **More flexible**
 - **Faster to solve: < 2 sec.**
 - **Given: 35 days / 20 PEDs / 7 shifts**

Metrics Formulation

- Feasibility problem
 - Constraint on metrics

$$\begin{aligned} \text{Min } & w_1(TSE) + w_2(NSE) + w_3(BSP) + w_4(PCC) \\ \text{s. t. } & \text{"rules/requirements"} \\ & x_{rds} \in \{0,1\} \\ & lb_{TSE} \leq (TSE) \leq ub_{TSE} \\ & lb_{NSE} \leq (NSE) \leq ub_{NSE} \\ & lb_{BSP} \leq (BSP) \leq ub_{BSP} \\ & lb_{PCC} \leq (PCC) \leq ub_{PCC} \end{aligned}$$

- Benefits of a feasibility problem
 - More flexible
 - Faster to solve: < 2 sec.
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Interactive Improvement

- Example output of metrics
 - Value (**Lower bound, Upper bound**)

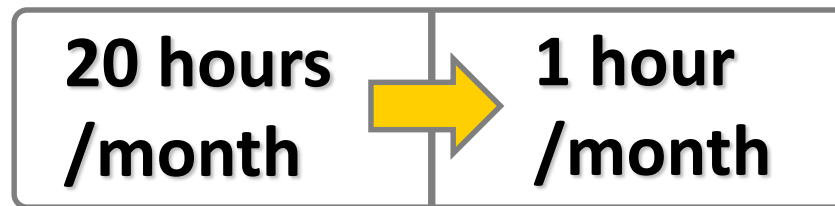
Resident Name	Number of Shifts	Number of Night Shifts	Number of Post CC	Number of Bad Sleep Templates
Smith	8 (7,9)	2 (0,10)	0 (0,1)	1 (0,1)
Sanchez	8 (7,10)	1 (0,10)	0 (0,1)	1 (0,1)
Chen	8 (7,9)	5 (0,10)	1 (0,1)	1 (0,1)
Shah	14 (13,15)	3 (0,10)	1 (0,1)	1 (0,1)
⋮	⋮	⋮	⋮	⋮

- Interactive approach engaging chief resident
 - Iteratively adjust bounds on metric constraints
 - Quickly build high quality schedule

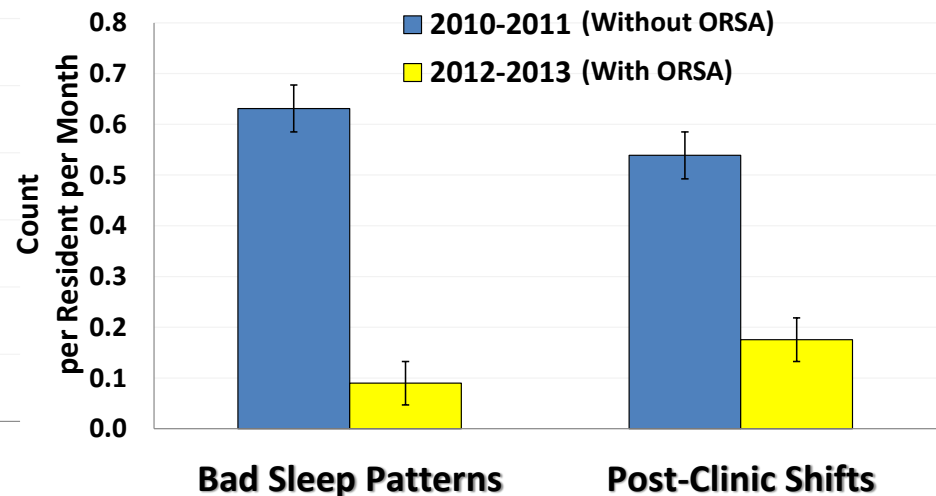
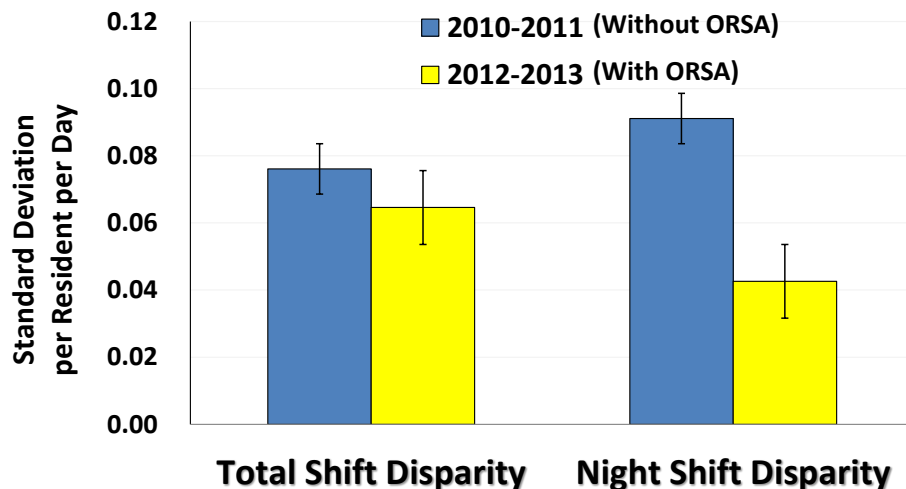
Results

- Our metrics-based scheduling tool:

- Reduces time to create schedules



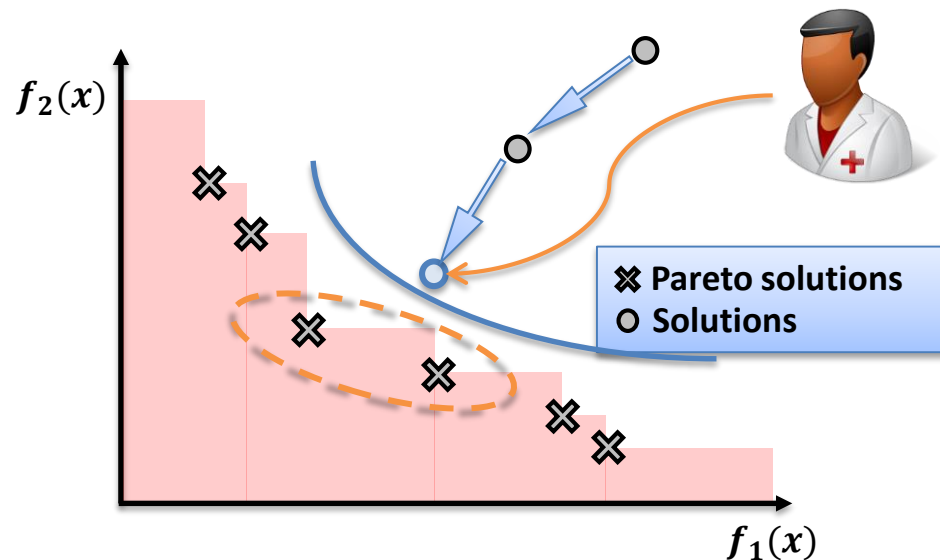
- Solves a multi-criteria scheduling problem



Limitations

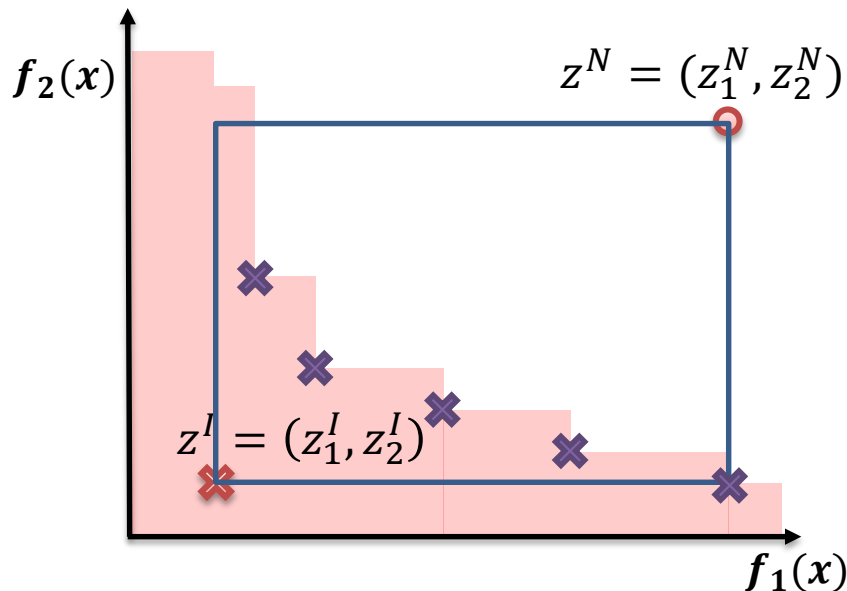
- **Myopic Solution**

- **Non-Pareto solution could be selected by chief residents**
 - Never see the whole picture (the set of Pareto solutions)
 - The most preferred solution is “most preferred” with respect to their satisfaction



Next Step

- Pareto Front
 - Generate the Pareto solutions of the problem (all of them or a sufficient representation)
 - Select the most preferred one among them



Efficient Schedules



Pareto: Bi-Objective Problem

- **Notation**

- \mathcal{H} : Solution Space, the set of feasible solutions
- \mathcal{P} : Pareto Front, the set of solutions in objective space
- $z_i = f_i(x)$: i th integer objective function, $\in \mathbb{Z}$
- Dominance ($<$) : $x < x'$ if and only if $z_i \leq z_i'$ where at least one inequality is strict

- **Bi-Objective Problem**

$$\begin{aligned} \min f(x) &= (f_1(x), f_2(x)) \\ \text{s. t. } x &\in \mathcal{H} \end{aligned}$$

Pareto: Bi-Objective Problem

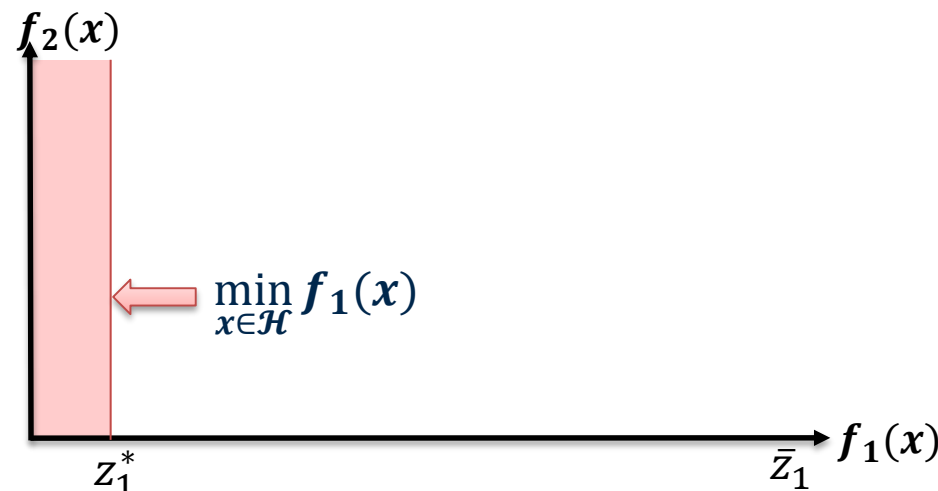
- Pareto Square Region

- Ideal Point:

- $z_1^* = \min_{x \in \mathcal{H}} f_1(x)$ and $z_2^* = \min_{x \in \mathcal{H}} f_2(x)$

- Nadir Point:

- $\bar{z}_1 = \min_{x \in \mathcal{H} \cap f_2(x)=z_2^*} f_1(x)$ and $\bar{z}_2 = \min_{x \in \mathcal{H} \cap f_1(x)=z_1^*} f_2(x)$



Pareto: Bi-Objective Problem

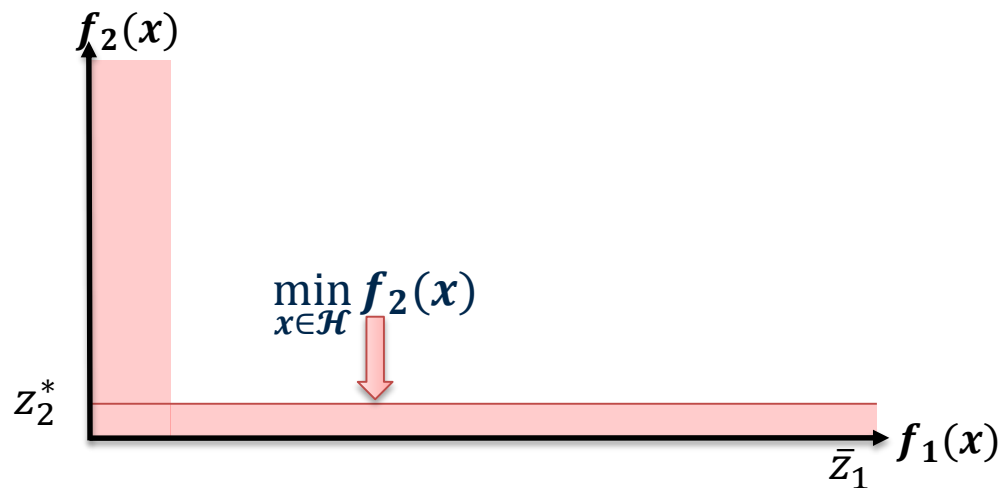
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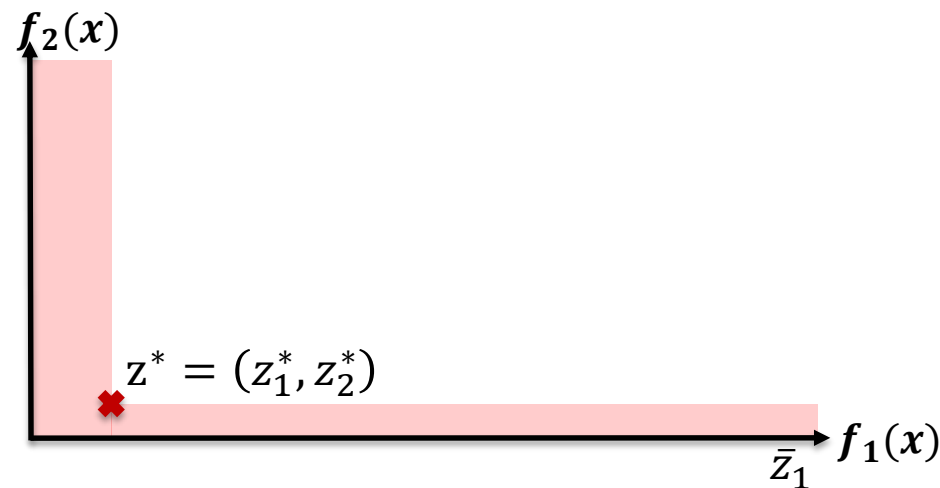
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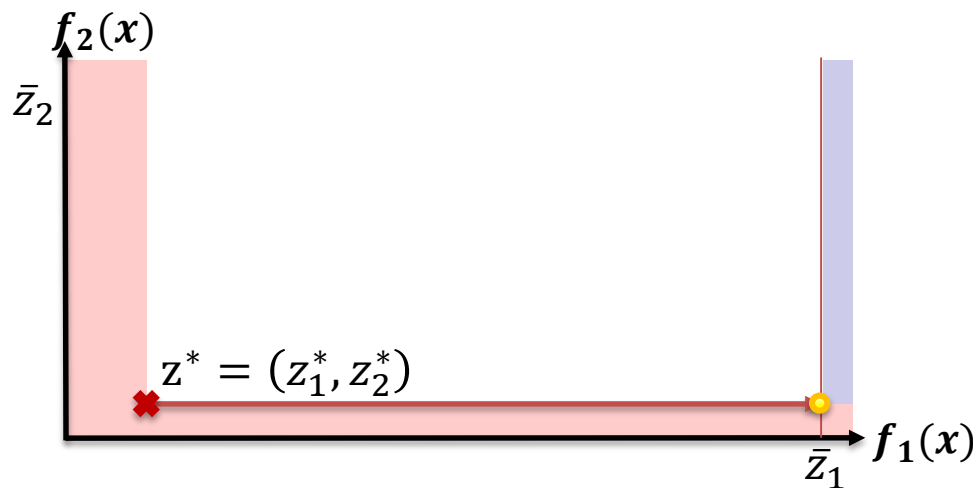
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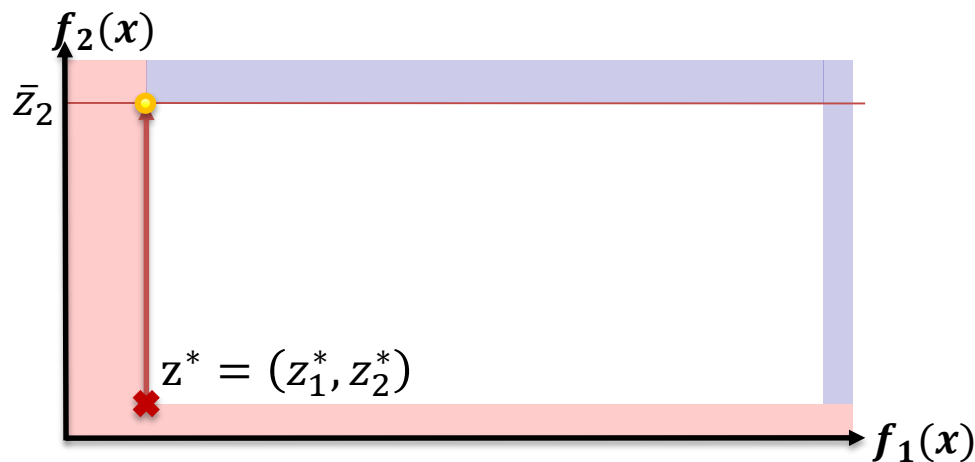
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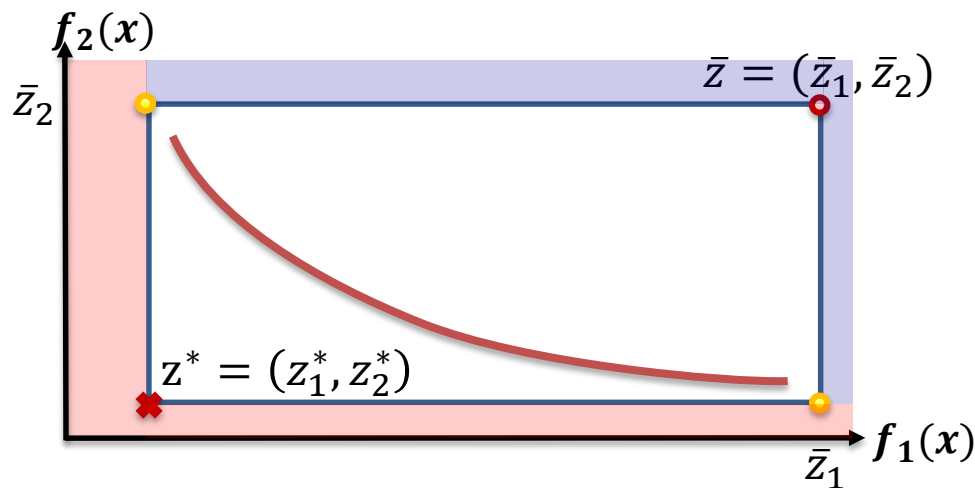
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Pareto: Bi-Objective Problem

• Pareto Squeezing Algorithm

Algorithm 1 Exact squeezing algorithm for bi-objective problems

Let P is set of pareto solutions we've found;

Compute the ideal (z_1^*, z_2^*) and Nadir (\bar{z}_1, \bar{z}_2) points;

Set $P := \{(\bar{z}_1, z_2^*)\}$ and $\delta := \bar{z}_1 - 1$;

WHILE $\delta \geq z_1^*$

 Solve $P_2(\delta)$ and get optimal solution (z_1^i, z_2^i) to $P_2(\delta)$;

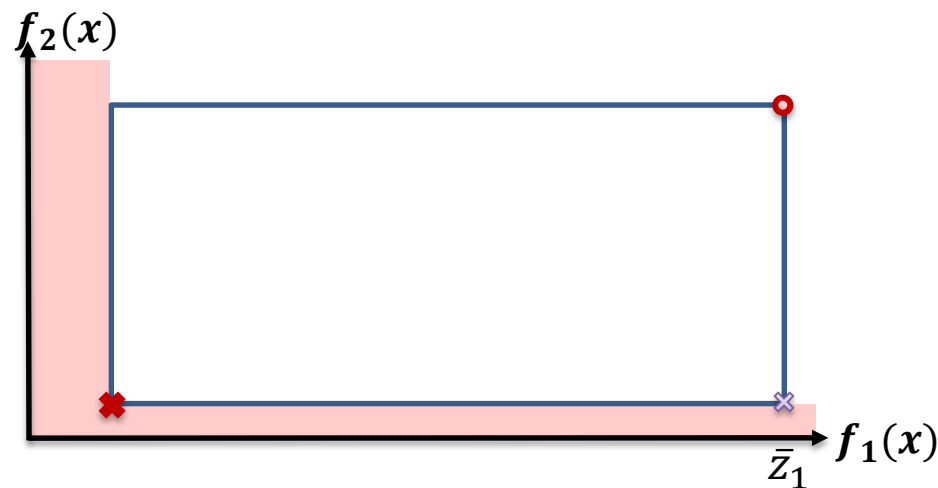
 //Given (z_1^i, z_2^i) , Find a left-bottom corner (\hat{z}_1^i) in the Pareto set;

 Solve $SQZ_1(z_2^i)$ and get optimal solution (\hat{z}_1^i, z_2^i) to $SQZ_1(z_2^i)$;

END WHILE

Set $P := P + (\hat{z}_1^i, z_2^i)$ and $\delta = \hat{z}_1^i - 1$;

GO Step 2 UNTIL $z_1 = z_1^*$;



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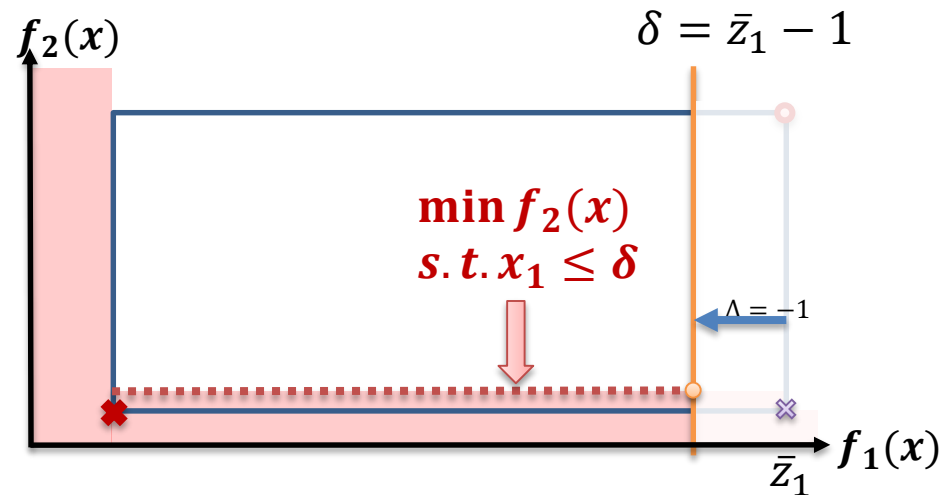
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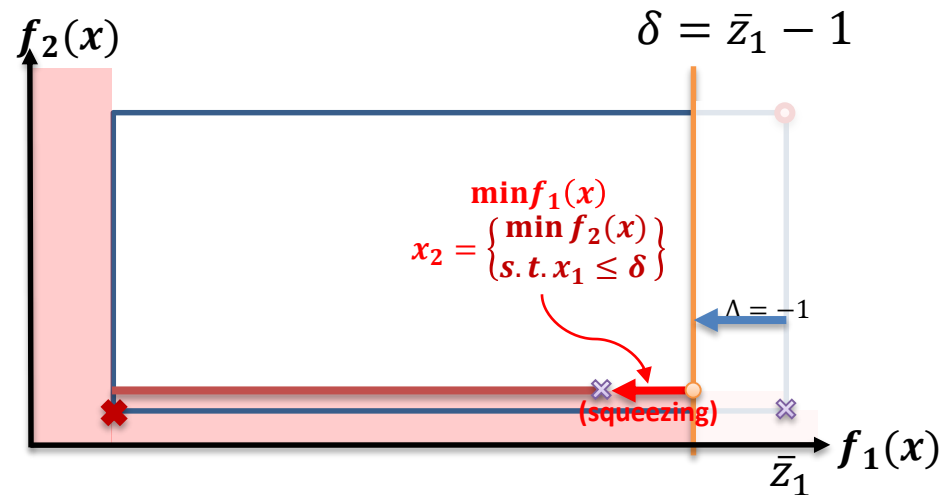
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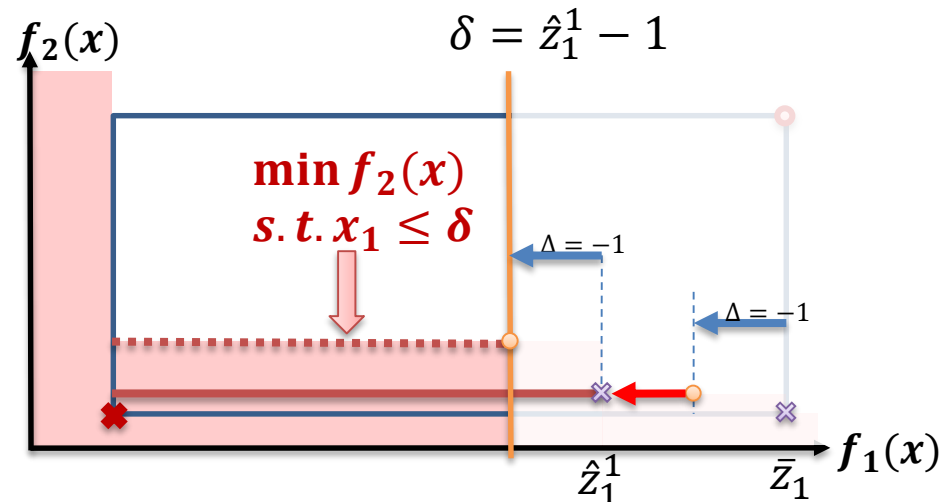
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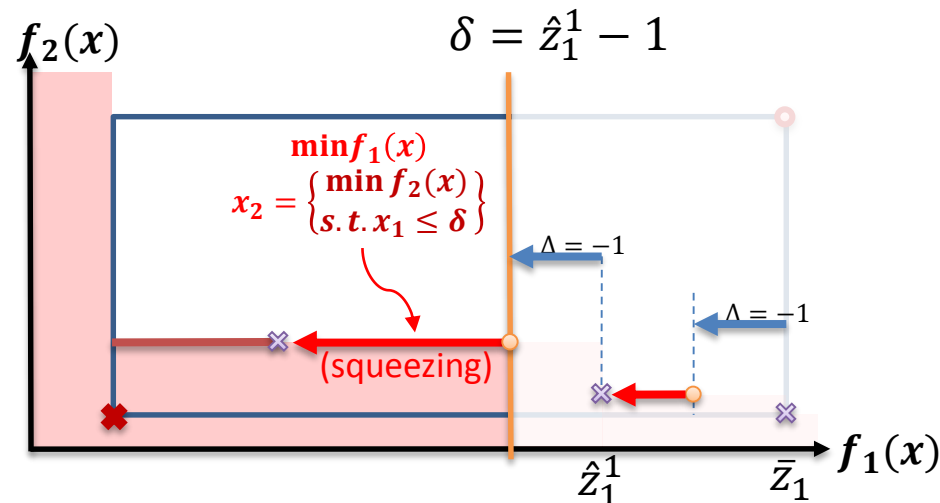
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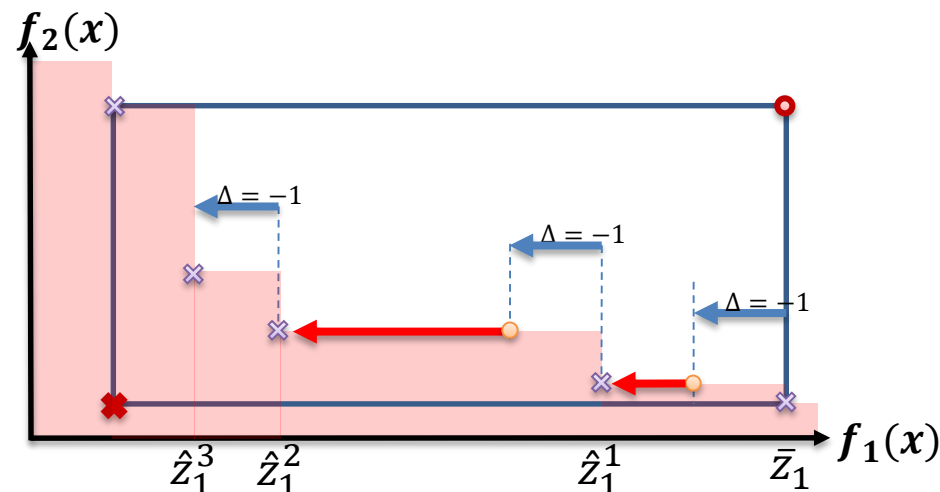
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Ongoing Research

- **Tri-Objective Problem**

$$\min f(x) = (f_1(x), f_2(x), f_3(x))$$
$$s. t. x \in \mathcal{H}$$

- **n -Objective Problem**

$$\min f(x) = (f_1(x), f_2(x), \dots, f_n(x))$$
$$s. t. x \in \mathcal{H}$$

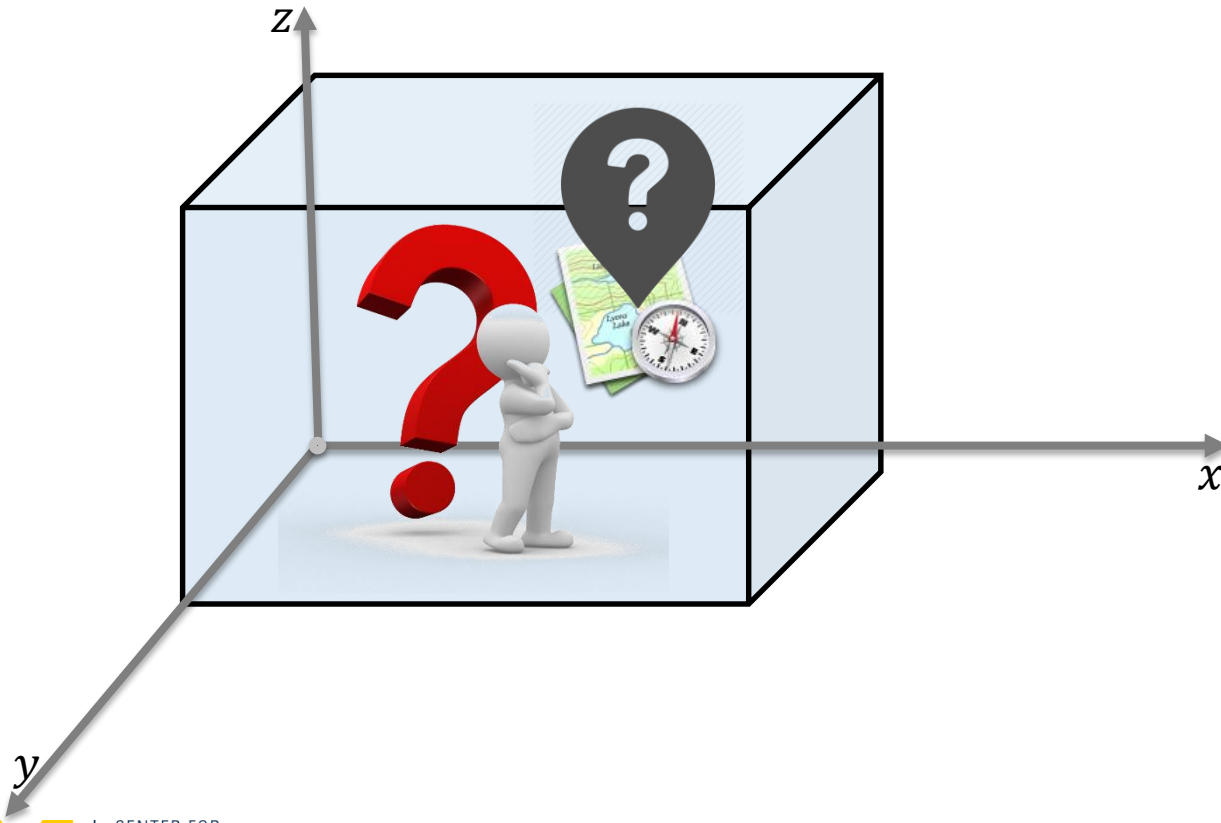
- **“Three is more than two plus one.”**

(Gandibleux, Xavier, ed. Multiple criteria optimization: state of the art annotated bibliographic surveys. Vol. 52. Springer Science & Business Media, 2002)



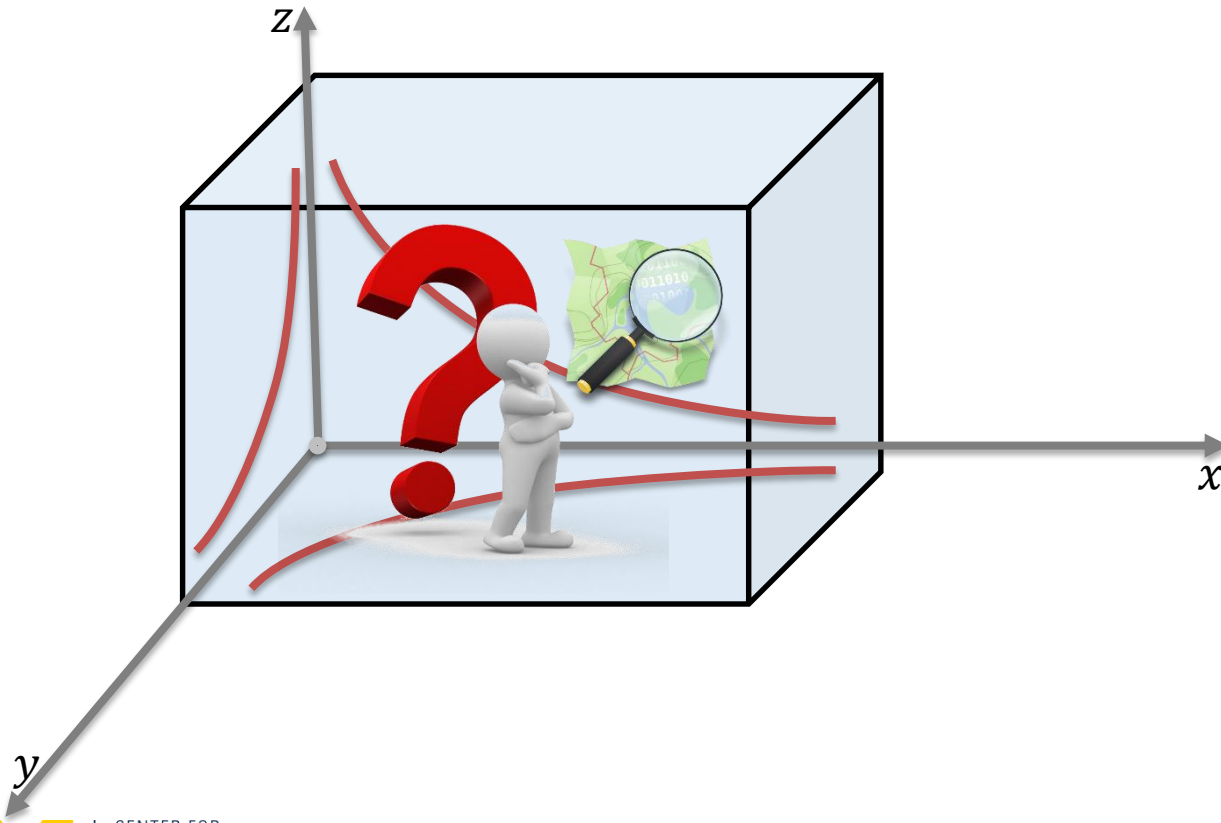
Pareto: Tri-Objective Problem

- **Tri-Objective Problem**
 - Where is Pareto Front?



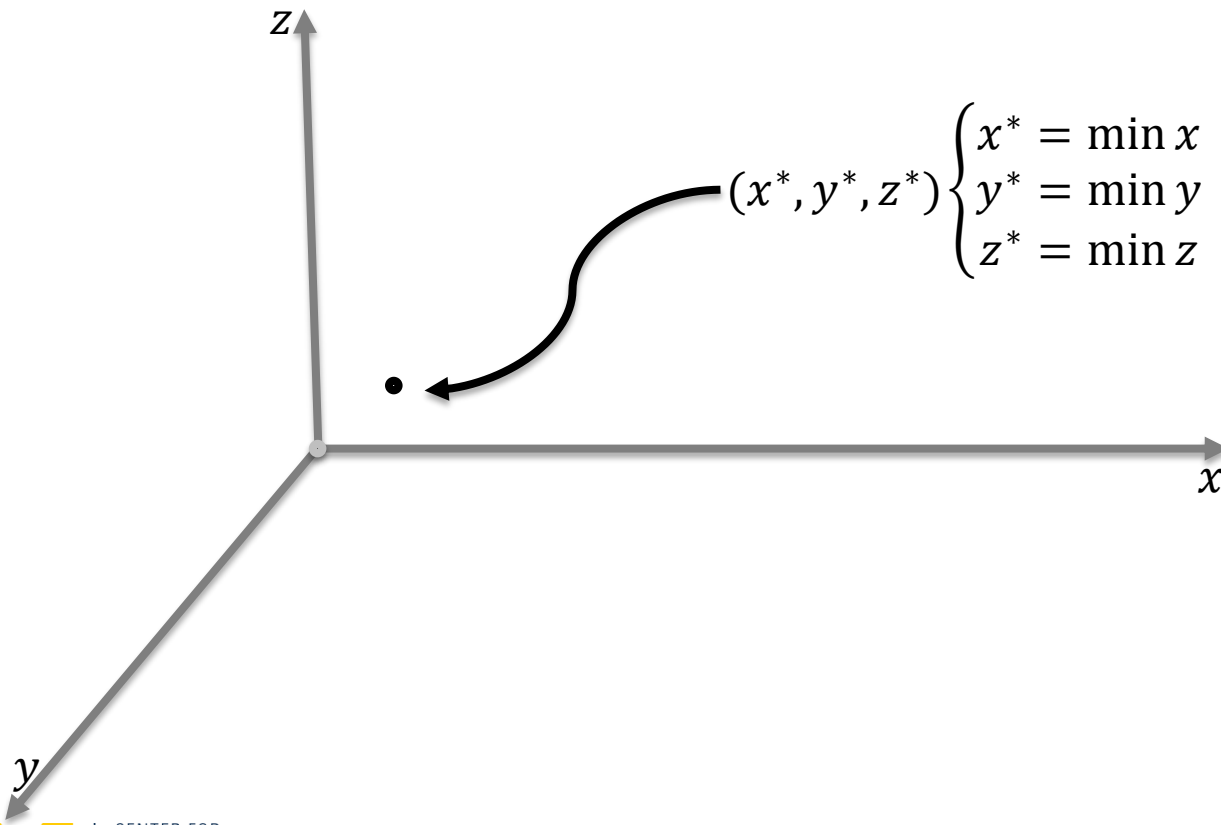
Pareto: Tri-Objective Problem

- **Tri-Objective Problem**
 - How to find Pareto Front?



Pareto: Tri-Objective Problem

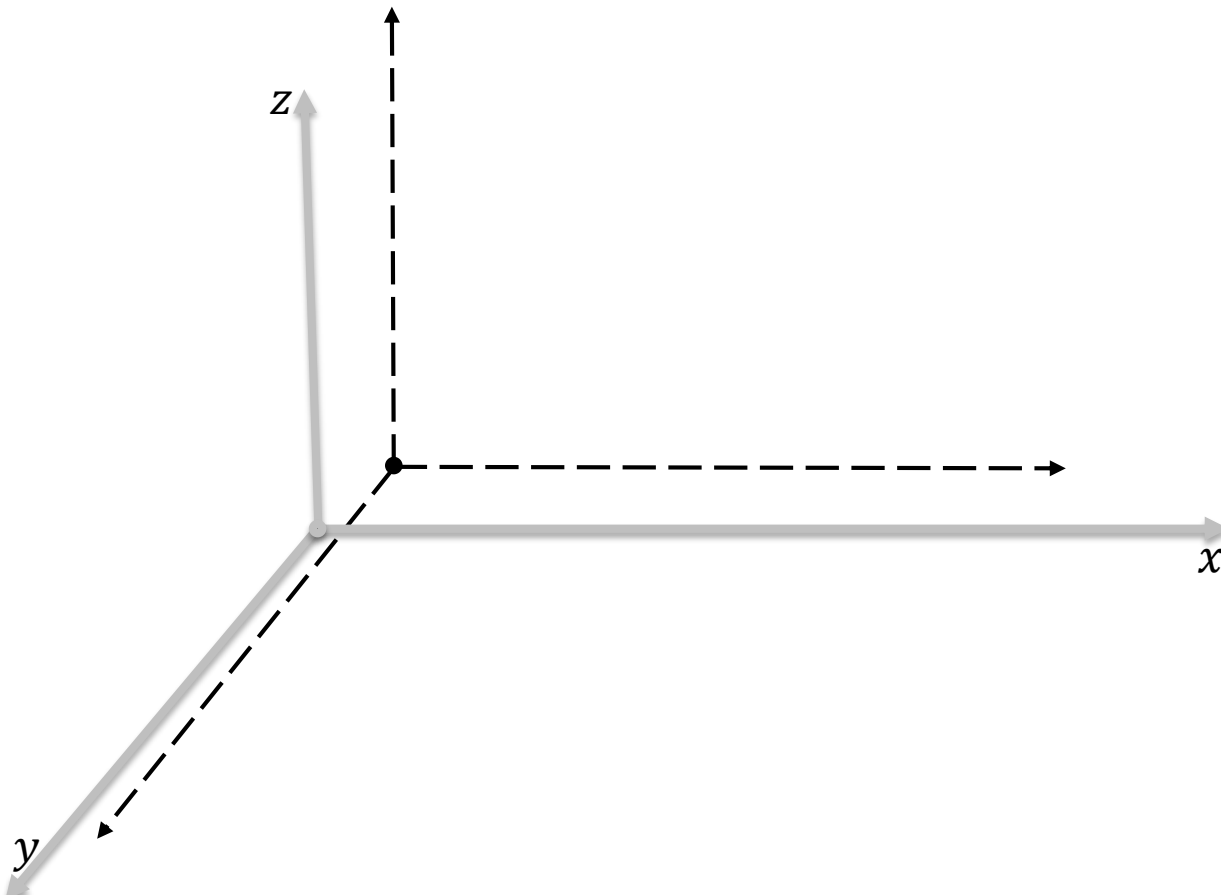
- Pareto Surface
 - Ideal Point and Nadir Points



Pareto: Tri-Objective Problem

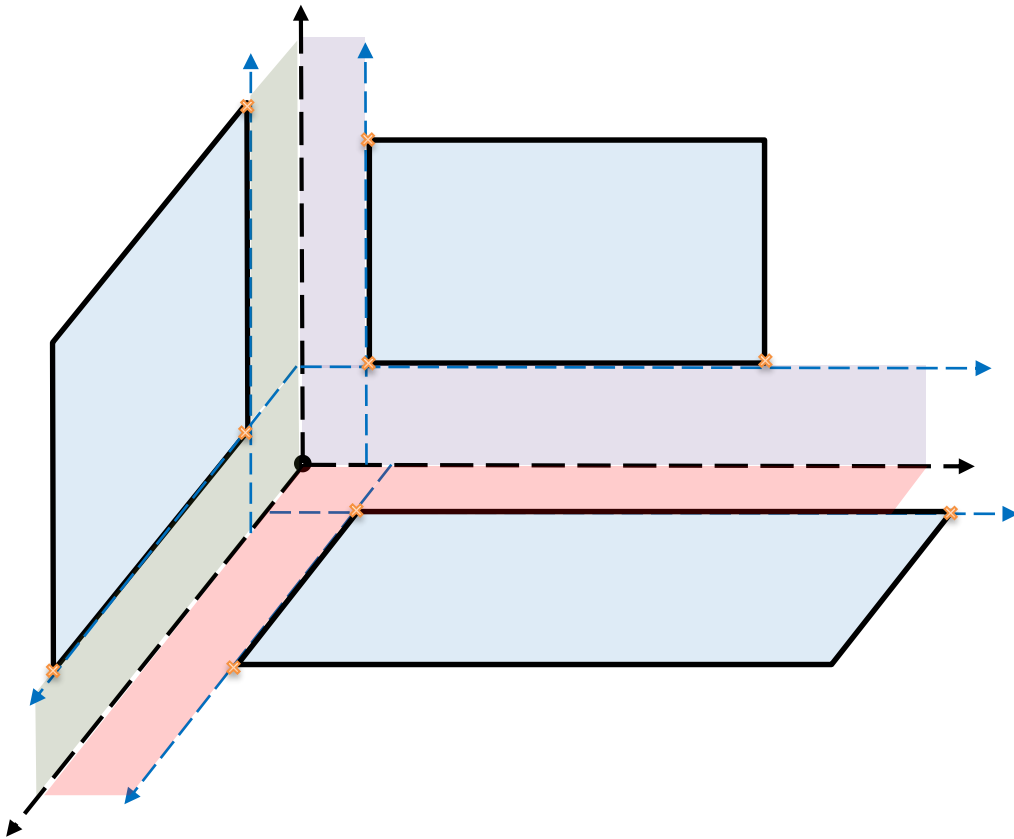
- Pareto Surface

- Ideal Point and Nadir Points



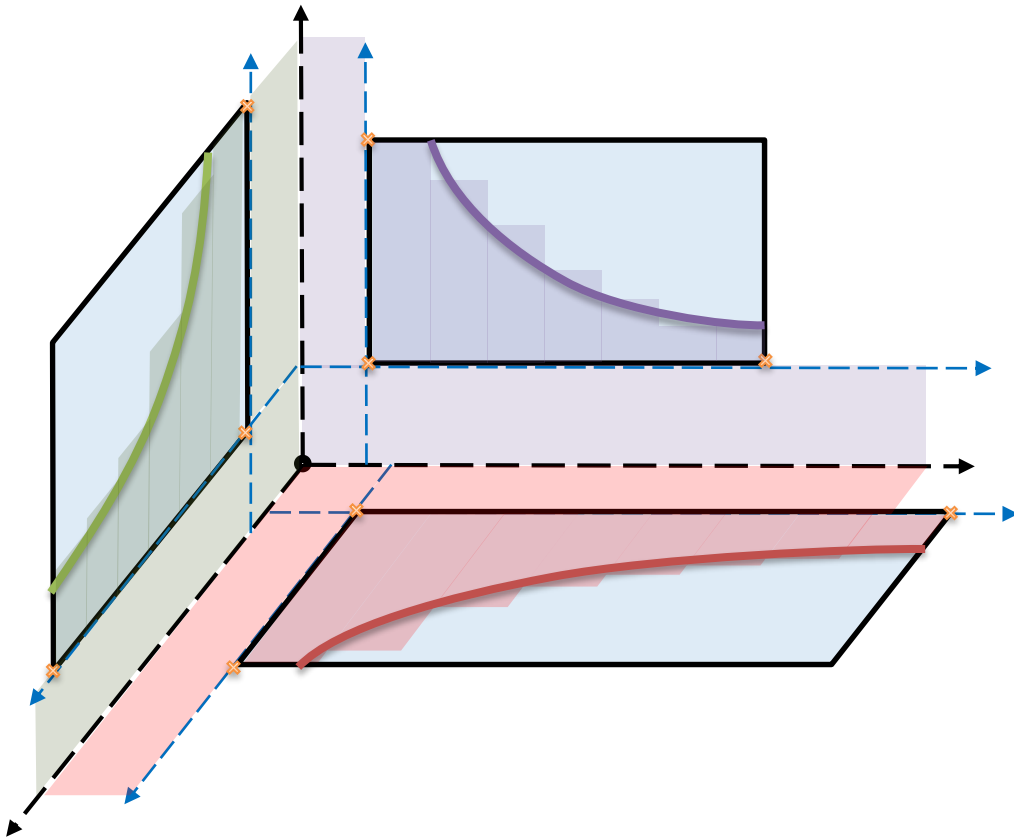
Pareto: Tri-Objective Problem

- Pareto Surface
 - Ideal Point and Nadir Points



Pareto: Tri-Objective Problem

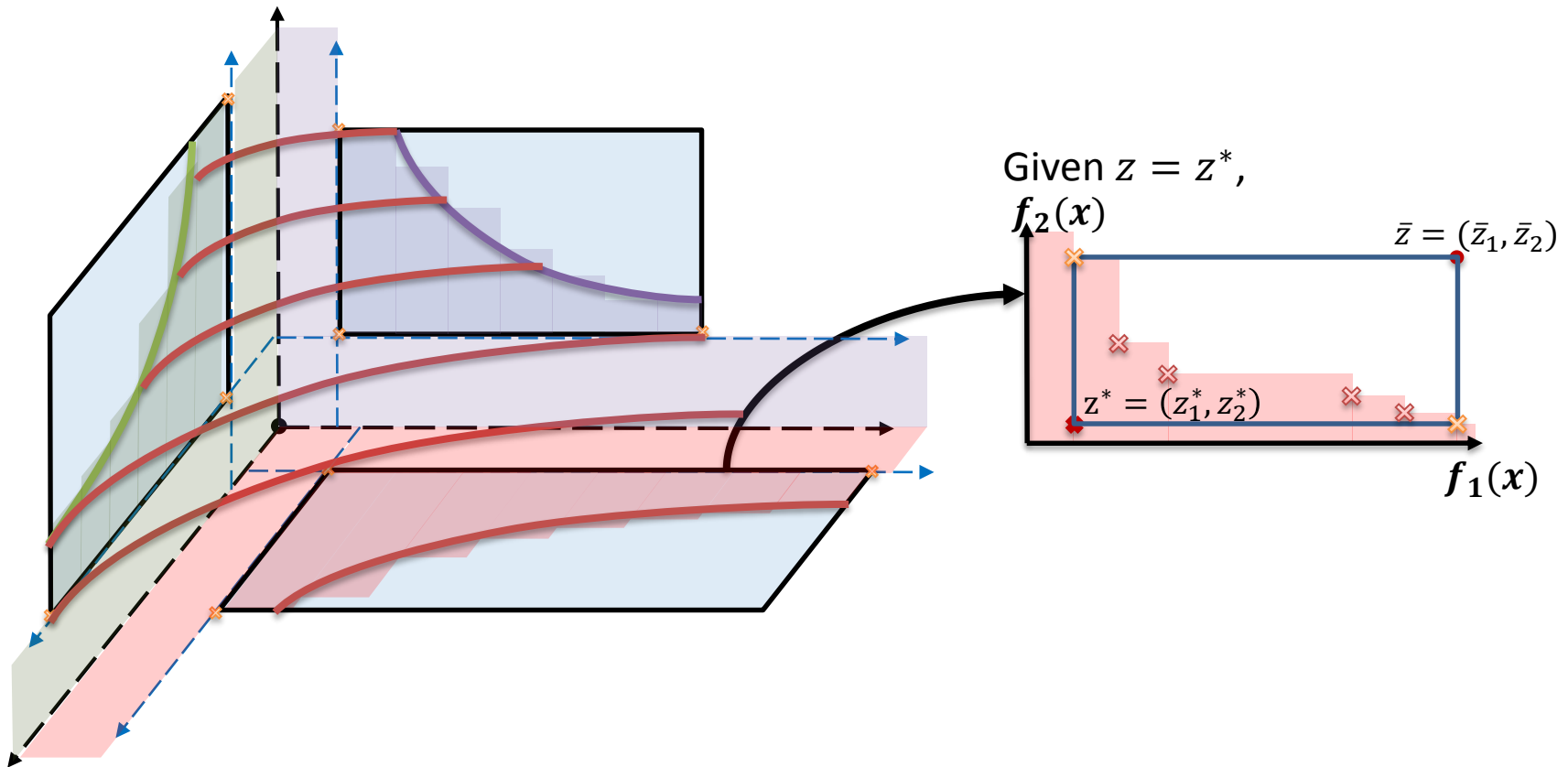
- Pareto Surface
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Pareto: Tri-Objective Problem

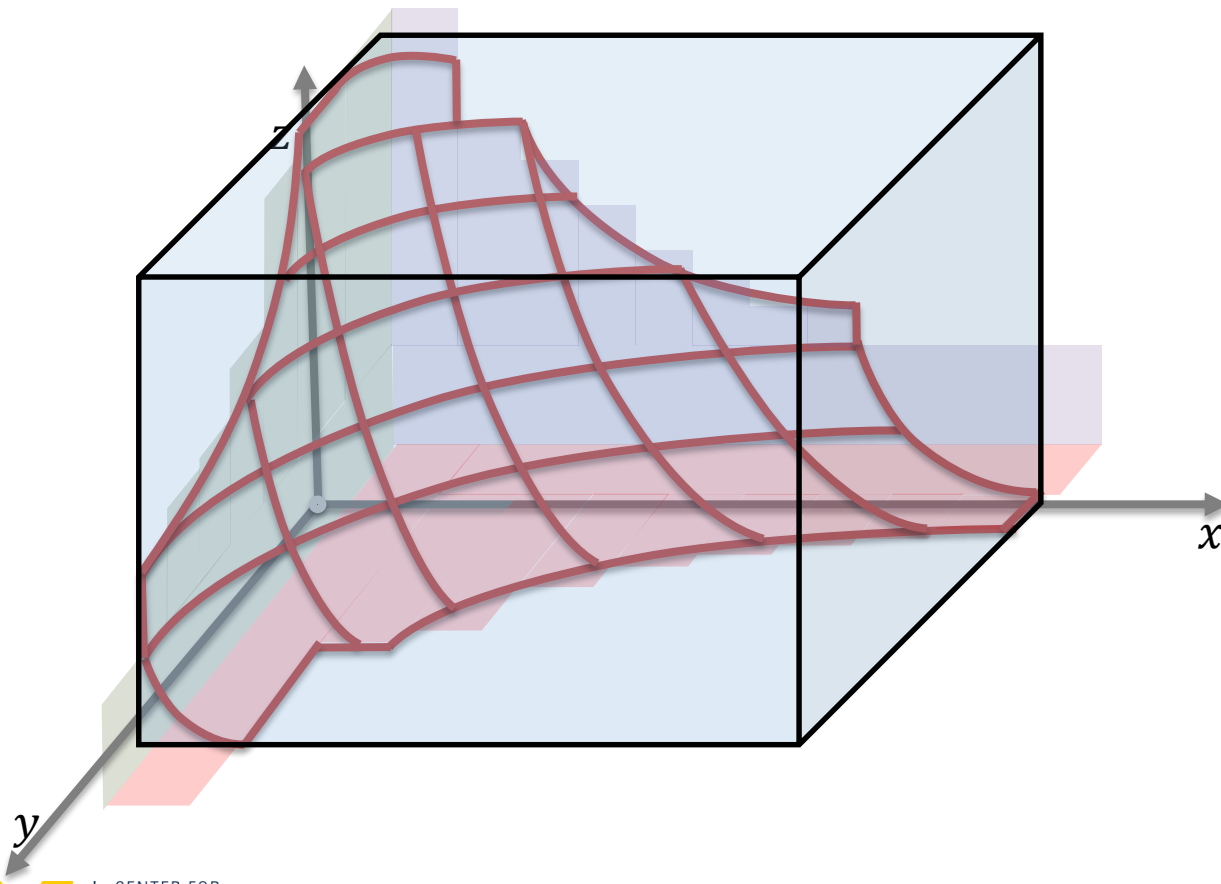
- Pareto Surface

- Given fixed z level, apply bi-objective approach



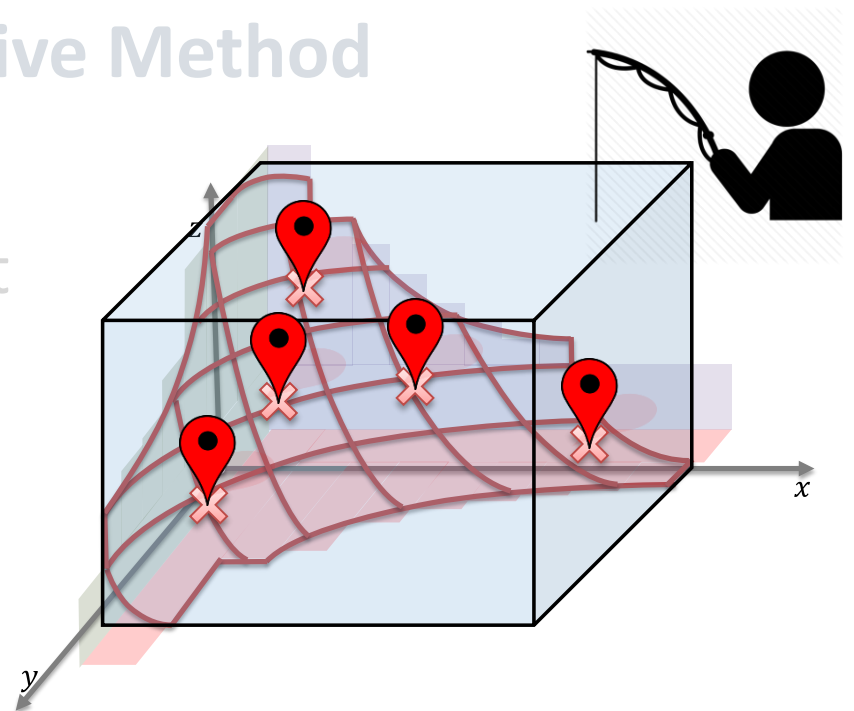
Pareto: Tri-Objective Problem

- Pareto Surface
 - Given fixed z level, apply bi-objective approach



Future Research

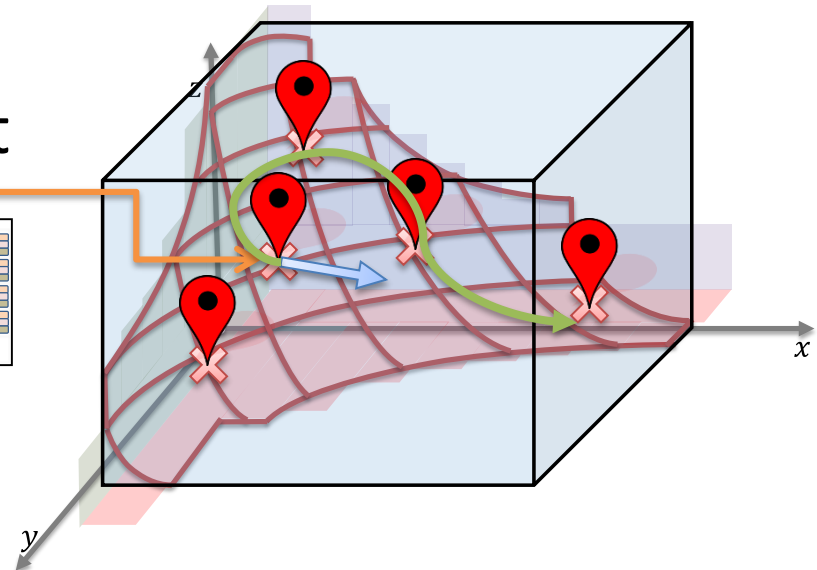
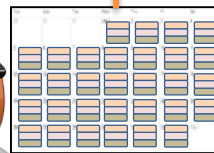
- **Approximation**
 - Sampling
 - Speed up
- **Integration with Interactive Method**
 - Interactive feedback
 - Decision on Pareto Front



Future Research

- **Approximation**
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Metrics	$f_1(x)$	$f_2(x)$	$f_3(x)$...
Δ Value	↑	↓	↑	...
Δ Value	↓	↑	↓	...
Δ Value	↓	↑	↓	...



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Feedback and Questions

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